Type-Theory of Parametric Algorithms

Roussanka Loukanova

Department of Algebra and Logic Institute of Mathematics and Informatics Bulgarian Academy of Sciences, Bulgaria

Online, 2020 Dec 18 Report Session — Отчетна Сесия за 2020 A glimpse at developments and applications, in my focus:

- Moschovakis [4], 1989: new theory of the mathematical notion of algorithm: untyped, full recursion
 - computations, by saving the algorithmic steps in memory locations
- ullet Moschovakis [5], 2006, Type-Theory of Acyclic Recursion, ${
 m L}_{
 m ar}^{\lambda}$

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- extended type systems variously, in Syn, Sem, SynSem; since 2007 in L^λ_{ar} & before in SitT
- extended reduction calculi in L_{ar}^{λ} , Loukanova, since 2011
- parametric / underspecified algorithms
- restrictor operator Loukanova [3], new 2020 & differently before
 constrained computations
 - restricted memory locations as generalised, restricted parameters
 - restricted objects
- Dependent-Type Theory of Full Recursion & Situated Information (DTTSitInfo)

This overview is about:

- ullet Type-Theory of (Acyclic) Algorithms, L^{λ}_{r} (L^{λ}_{ar}): a new approach to
 - the math notion of algorithm
 - Computational Semantics CompSem (AL / FL /NL / HL)

Placement of $L_{\mathrm{ar}}^{\lambda}$ in a class of type theories

Montague IL
$$\subsetneq$$
 Gallin TY $_2 \subsetneq$ Moschovakis $L_{ar}^{\lambda} \subsetneq$ Moschovakis L_r^{λ} (1)

$$\stackrel{?}{\equiv} \mathsf{DTTSitInfo} \tag{2}$$

- Montague IL for PTQ (1970-73)
- Gallin TY₂ (1975)
- Acyclic (closing-off) Recursion, L^λ_{ar}
 Moschovakis [5] (2006), Loukanova (since 2007)
- ullet L $_{
 m ar}^{\lambda}$ is a higher-order type theory with terms for acyclic recursion:
 - \bullet the algorithmic meanings of the meaningful L_{ar}^{λ} terms are acyclic algorithms that close-off
- \bullet $L^{\lambda}_{\rm r}$ is a higher-order type theory with full recursion
- Dependent-Type Theory of Full Recursion & Situated Information (DTTSitInfo)

Some Classes of Type Theories (TTs)

$$\mathsf{IL} \subsetneq \mathsf{TY}_2 \subsetneq \cdots \subsetneq \mathsf{L}^{\lambda}_{\mathrm{ar}} \subsetneq \mathsf{L}^{\lambda}_{\mathrm{ra}} \subsetneq \mathsf{L}^{\mathrm{ta}}_{\mathrm{ar}} \subsetneq \mathsf{L}^{\lambda}_{\mathrm{r}} \subsetneq \mathsf{L}^{\mathrm{ta}}_{\mathrm{r}}$$
 (3)

$$L_{ar}^{\lambda} \subsetneq L_{ra}^{\lambda} \subsetneq L_{ar}^{ta} \subsetneq L_{r}^{ta} \subsetneq \mathsf{DTTSitInfo} \stackrel{?}{\subsetneq} L_{gp}^{st} \tag{4}$$

- Montague IL for PTQ (1970-73) | Gallin TY₂ (1975)
- \bullet L_{ar}^{λ} , L_{ra}^{λ} , L_{r}^{λ} , and L_{ar}^{ta} , L_{r}^{ta} are TTs of Algorithms (2006–ongoing)
 - Recursion Operator: Recursion Terms, Moschovakis, 1989

Reduction Calculi in the $L_{\mathrm{ar}}^{\lambda}$ Classes via Reduction Rules

- ullet Reducing every term A to a canonical form $\operatorname{cf}(A)$
- \bullet cf (A) of an algorithmically meaningful A determines the algorithm ${\rm alg}(A)$ for computing ${\rm den}(A),$ in a step-by-step mode
 - 2 Restrictor Operator: Restrictor Terms, Loukanova (2020; earlier vars)
 - \bullet L_{ar}^{ta} , L_{r}^{ta} polymorphic types
 - **Q** Reduction Calculi, for each extended L_{ar}^{λ} η -reduction, η^* -reduction (in prog.), γ -reduction, γ^* -reduction

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Gallin Types:
                                                   \sigma :\equiv e \mid t \mid s \mid (\tau_1 \rightarrow \tau_2) (Gallin, 1975)
   For all \tau \in \mathsf{Types}:
   Constants:
                                                   \mathsf{Consts}_{\tau} = \{\mathsf{c}_0^{\tau}, \mathsf{c}_1^{\tau}, \dots, \mathsf{c}_k^{\tau}\}
   Variables:
                                                   PureV_{\tau} = \{v_0^{\tau}, v_1^{\tau}, \dots\}
                                                    \mathsf{MemoryV}_{\tau} = \mathsf{RecV}_{\tau} = \{p_0^{\tau}, p_1^{\tau}, \dots\}
Terms of L_{ar}^{\lambda} (L_{r}^{\lambda}):
   A :\equiv c^{\tau} : \tau \mid x^{\tau} : \tau \quad \text{(for } c^{\tau} \in \mathsf{Consts}_{\tau}, \ x^{\tau} \in \mathsf{PureV}_{\tau} \cup \ \mathsf{RecV}_{\tau} \text{)}
                                                                                                                                                                 (5a)
                               |\mathsf{B}^{(\sigma \to \tau)}(\mathsf{C}^{\sigma}) : \tau
                                                                                                                                                                  (5b)
                               |\lambda(v^{\sigma})(\mathsf{B}^{\tau}):(\sigma\to\tau) \quad \text{(for } v^{\sigma}\in\mathsf{PureV}_{\sigma}\text{)}
                                                                                                                                                                 (5c)
                               | A_0^{\sigma_0}  where | p_1^{\sigma_1} := A_1^{\sigma_1}, \ldots, 
                                                                                                                                                                 (5d)
                                                                   p_i^{\sigma_i} := A_i^{\sigma_i}, \dots, p_n^{\sigma_n} := A_n^{\sigma_n} \} ] : \sigma_0
                               \left[\mathsf{A}_0^{\sigma_0} \text{ such that } \left\{ \mathsf{C}_1^{\tau_1}, \ldots, \mathsf{C}_m^{\tau_m} \right\} \right] : \sigma_0'
                                                                                                                                                                  (5e)
     • B, C \in Terms, p_i^{\sigma_i} \in \text{RecV}_{\sigma_i}, A_i^{\sigma_i} \in \text{Terms}_{\sigma_i}
          C_i^{\tau_j} \in \mathsf{Terms}_{\tau_i} (for propositions): \tau_i \equiv \mathsf{t} or \tau_i \equiv \widetilde{\mathsf{t}} \equiv (\mathsf{s} \to \mathsf{t})
     • Acyclicity Constraint, for L_{n}^{\lambda}; without it, L_{n}^{\lambda}
           \{p_1^{\sigma_1} := \mathsf{A}_1^{\sigma_1}, \dots, p_i^{\sigma_i} := \mathsf{A}_i^{\sigma_i}, \dots, p_n^{\sigma_n} := \mathsf{A}_n^{\sigma_n}\} is acyclic iff:
                 • there is a rank: \{p_1, \ldots, p_n\} \to \mathbb{N} such that:
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if $p_i \in \mathsf{FreeVars}(A_i)$ then $\mathsf{rank}(p_i) > \mathsf{rank}(p_i)$

Outlook1: Development of Computational Theories and Applications

- Generalised Computational Grammar:
 via faithfull representations of CompSynSem interfaces in HL
 - Hierarchical lexicon with morphological structure and lexical rules
 - Lexicon that propagates into the phrasal structure of sentences
 - Syntax of HL expressions (phrasal and grammatical dependences)
 - Syntax-semantics inter-relations in lexicon and phrases
 - Abstract and Parametric Grs: across syn categories and semantic types vs.

Specific Instantiations of lexemes and phrases: in and across languages

• A Big Picture — simplified and approximated, but realistic:

(I've done quite a lot of it, but still a lot to do!)

Some Current Tasks (among many others) and Future Work

- My focus is on:
 - \bullet Development of L_{ar}^{λ} and L_{r}^{λ}
 - Dependent-Type Theory of Situated Information and Algorithms
 - Applications to formal and natural languages
 - Extending the Coverage of Computational Semantics
 - Computational Syntax-Semantics Interfaces
 - Semantics of programming and specification languages
 - Theoretical foundations of (parts of) compilers
- More to come

THANK YOU!

Some References I



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