

<p>Provided for non-commercial research and educational use. Not for reproduction, distribution or commercial use.</p>
--

# Mathematica Balkanica

Mathematical Society of South-Eastern Europe  
A quarterly published by  
the Bulgarian Academy of Sciences – National Committee for Mathematics

---

The attached copy is furnished for non-commercial research and education use only.  
Authors are permitted to post this version of the article to their personal websites or  
institutional repositories and to share with other researchers in the form of electronic  
reprints.

Other uses, including reproduction and distribution, or selling or licensing copies, or  
posting to third party websites are prohibited.

For further information on Mathematica Balkanica visit the website of the journal  
<http://www.mathbalkanica.info>

or contact:

Mathematica Balkanica - Editorial Office;  
Acad. G. Bonchev str., Bl. 25A, 1113 Sofia, Bulgaria  
Phone: +359-2-979-6311, Fax: +359-2-870-7273,  
E-mail: [balmat@bas.bg](mailto:balmat@bas.bg)

# Boolean Ring Equations <sup>1</sup>

*Dragić Banković*

*Presented by Ž. Mijajlović*

If a particular solution of an equation over a Boolean ring is known, Löwenheim's theorem determines a general solution of this equation (see for instance, [7]). In this paper we reduce the finding of a particular solution to solving simple equations in two-element Boolean ring, i.e. on  $\{0,1\}$ . In the similar way we also determine general solutions of any equations over a Boolean ring, where we do not suppose that a particular solution is known.

## 1. Introduction

Solving of equations is a basic inference mechanism in algebraic manipulation of formulas, automated reasoning and some programming languages. Since this paper is concerning with general solution we firstly state the definition of a general solutions.

**Definition 1.** Let  $E$  be a given non-empty set and  $Q$  be a given unary relation of  $E$ . A formula  $x = \varphi(t)$ , where  $\varphi: E \rightarrow E$  is a given function, represents a general solution of the  $x$ -equation  $Q(x)$  if and only if

$$(\forall t)Q(\varphi(t)) \wedge (\forall x)(Q(x) \Rightarrow (\exists t)x = \varphi(t)).$$

We say that the equation  $Q(x)$  is consistent if there is  $y \in E$  such that  $Q(y)$  is true.

## 2. Boolean functions and Boolean polynomials

Let  $N = \{1, \dots, n\}$ , where  $n$  is a natural number.

<sup>1</sup>This research was supported by Science Found of Serbia, Grant number 0401A, through Matematički institut

**Theorem 1 [9].** *A mapping  $f: B^n \rightarrow B$  is a Boolean function if and only if it can be written in the canonical disjunctive form*

$$f(X) = \bigcup_A f(A)X^A,$$

where  $\bigcup_A$  means the union over all  $A \in \{0, 1\}^n$ .  $f$  is a simple Boolean function if  $(\forall A \in \{0, 1\}^n) f(A) \in \{0, 1\}$ .

**Theorem 2 [9].** *A mapping  $f: B^n \rightarrow B$  is a Boolean polynomial if and only if it can be written in the canonical polynomial form*

$$f(X) = \sum_{s \subset N} b_s \prod_{i \in s} x_i.$$

$f$  is a simple Boolean polynomial if and only if  $(\forall s \in N) b_s \in \{0, 1\}$ .

**Theorem 3 [9].** *Let  $\mathcal{B} = (B, \cup, \cdot, ', 0, 1)$  be a Boolean algebra and  $n$  a natural number. A mapping  $f: B^n \rightarrow B$  is a Boolean function if and only if it is polynomial of the Boolean ring  $\mathcal{R} = (B, +, \cdot, 0, 1)$ .*

### 3. Boolean equations

In this section we collect some background material about Boolean equations. A detailed account is given in [9]. We shall use the notation:  $X = (x_1, \dots, x_n)$  and  $T = (t_1, \dots, t_n)$ .

Given an arbitrary Boolean algebra  $\mathcal{B}$ , a Boolean equation in  $n$  unknowns over  $\mathcal{B}$  is an equation of the form  $g(X) = h(X)$ , where  $g, h: B^n \rightarrow B$  are Boolean functions.

**Theorem 4 [9].** *Every Boolean equation or system of Boolean equations is equivalent to a single Boolean equation of the form  $f(X) = 0$ , where  $f$  is a Boolean function.*

**Theorem 5 [9].** *The Boolean equation  $f(X) = 0$  is consistent if and only if*

$$\prod_A f(A) = 0,$$

where  $\prod_A$  means the product over all  $A \in \{0, 1\}^n$ .

**Theorem 6 [4].** *Let  $f, g_1, \dots, g_n: B^n \rightarrow B$  be Boolean functions and  $G = (g_1, \dots, g_n)$ . The formula*

$$X = G(T)$$

(or, in scalar form,  $x_j = g_j(t_1, \dots, t_n)$  ( $j = 1, \dots, n$ )) represents a general solution of the consistent equation  $f(X) = 0$  if and only if

$$(\forall T)(f(T) = \prod_A \bigcup_{j=1}^n (g_j(A) + t_j)).$$

**Theorem 7 [9].** Let  $f, g: B^n \rightarrow B$  be Boolean functions and assume that  $\prod_A f(A) = 0$ . Then the following conditions are equivalent:

$$\begin{aligned} &(\forall X \in B^n)(f(X) = 0 \Rightarrow g(X) = 0) \\ &(\forall X \in B^n)(g(X) \leq f(X)) \\ &(\forall X \in \{0, 1\}^n)(g(X) \leq f(X)). \end{aligned}$$

We prove now a lemma that will be used in the proofs of Theorem 9 and Theorem 10.

**Lemma.** Let  $g: B^n \rightarrow B$  be a Boolean function and  $f: B^n \rightarrow B$  be a simple Boolean function. If  $\prod_A f(A) = 0$ , then the conditions

- (a)  $(\forall X \in B^n)(f(X) = 0 \Rightarrow g(X) = 0)$
- (b)  $(\forall A \in \{0, 1\}^n)(f(A) = 0 \Rightarrow g(A) = 0)$

are equivalent.

**Proof.** (a)  $\rightarrow$  (b) is trivial. Since  $f$  is the simple Boolean function, we have  $(\forall A \in \{0, 1\}^n)(f(A) \in \{0, 1\})$ . Let (b) hold. If  $f(A) = 0$ , then  $g(A) = 0$ , i.e.  $g(A) \leq f(A)$ . If  $f(A) = 1$ , then  $g(A) \leq f(A) = 1$ . Therefore we have  $(\forall A \in \{0, 1\}^n)(g(A) \leq f(A))$ . Since the latter formula is equivalent to (a), by Theorem 3, (b)  $\rightarrow$  (a) is proved. ■

#### 4. Boolean ring equations

Given an arbitrary Boolean ring  $\mathcal{R}$ , a Boolean ring equation in  $n$  unknowns over  $\mathcal{B}$  is an equation of the form  $g(X) = h(X)$ , where  $g, h: B \times B \rightarrow B$  are Boolean polynomials. Taking in mind Theorem 3 and Theorem 4, every Boolean ring equation is equivalent to a single Boolean ring equation of the form  $f(X) = 0$ , where  $f$  is a Boolean polynomial.

**Theorem 8 [9].** The Boolean ring equation

$$\sum_{s \in N} b_s \prod_{i \in S} x_i = 0$$

is consistent (i.e.  $(\exists X) \sum_{s \in N} b_s \prod_{i \in S} x_i = 0$ ) if and only if

$$b_{\emptyset} \prod_{\emptyset \neq S \subseteq N} (b_s + 1) = 0.$$

#### 4.1 Particular solutions

**Theorem 9.** Let  $Y$  is the  $m$ -tuple ( $m < 2^n$ ) of all different elements  $b_s$  from  $\sum_{s \in N} b_s \prod_{i \in S} x_i = 0$  and let  $h(X, Y) = \sum_{s \in N} b_s \prod_{i \in S} x_i$ . If  $b_{\emptyset} \prod_{\emptyset \neq S \subseteq N} (b_s + 1) = 0$ , then the formulas

$$(1) \quad p_j = \sum_C z_{j,C} Y^C \quad (j = 1, \dots, n)$$

( $\sum_C$  means the sum over all  $C \in \{0, 1\}^m$ ) represents a particular solution of Boolean ring equation  $h(X, Y) = 0$  with respect to  $X$ , if and only if

$$(2) \quad (\forall Y \in V) h(z_{1,Y}, \dots, z_{n,Y}, Y) = 0,$$

where

$$V = \{Y | Y \in \{0, 1\}^m \wedge (\exists X) h(X, Y) = 0\}, \text{ i.e.}$$

$$(3) \quad V = \{Y | Y \in \{0, 1\}^m \wedge b_{\emptyset} \prod_{\emptyset \neq S \subseteq N} (b_s + 1) = 0\}.$$

**Proof.** It is obvious that a particular solution  $(p_1, \dots, p_n)$  of the equation  $h(X, Y) = 0$  depends on  $Y$ , i.e. it is of the form

$$p_j = \sum_C z_{j,C} Y^C \quad (j = 1, \dots, n),$$

because of Theorem 3 and Lemma.

Therefore we have the following equivalences:

$$(\forall Y \in B^m) (b_{\emptyset} \prod_{\emptyset \neq S \subseteq N} (b_s + 1) = 0 \Rightarrow h(\sum_C z_{1,C} Y^C, \dots, \sum_C z_{n,C} Y^C, Y) = 0)$$

$$\begin{aligned} &\Leftrightarrow (\forall Y \in \{0, 1\}^m) (b_{\emptyset} \prod_{\emptyset \neq S \subseteq N} (b_s + 1) = 0 \\ &\Rightarrow h(\sum_C z_{1,C} Y^C, \dots, \sum_C z_{n,C} Y^C, Y) = 0) \end{aligned}$$

(because of Theorem 4 and Lemma)

$$\Leftrightarrow (\forall Y \in V) h(\sum_C z_{1,C} Y^C, \dots, \sum_C z_{n,C} Y^C, Y) = 0$$

$$(V = \{Y | Y \in \{0, 1\}^m \quad \wedge \quad b_\emptyset \prod_{\emptyset \neq S \subseteq N} (b_s + 1) = 0\})$$

$$\Leftrightarrow (\forall Y \in V) h(z_{1,Y}, \dots, z_{n,Y}, Y) = 0$$

(because  $(\forall Y \in V)(\forall j \in N)(\sum_C z_{j,C} Y^C = z_{j,Y})$ . ■

The algorithm for solving the latter system has been given in [7].

If we find  $z_{1,Y}, \dots, z_{n,Y}$  we get the particular solution (1).

#### 4.2 General solutions

**Theorem 10.** *Let  $Y$  be the  $m$ -tuple ( $m < 2^n$ ) of all different elements  $b_s$  from  $\sum_{S \subseteq N} b_S \prod_{i \in S} x_i = 0$  and let  $h(X, Y) = \sum_{S \subseteq N} b_S \prod_{i \in S} x_i$ . If  $b_\emptyset \prod_{\emptyset \neq S \subseteq N} (b_s + 1) = 0$ , then the formulas*

$$(4) \quad x_j = \sum_D \left( \sum_C u_{j,D,C} Y^C \right) T^D \quad (j = 1, \dots, n)$$

( $\sum_D$  means the sum over all  $D \in \{0, 1\}^n$  and  $\sum_C$  means the sum over all  $C \in \{0, 1\}^m$ ) represent a general solution of Boolean ring equation  $h(X, Y) = 0$ , with respect to  $X$ , if and only if

$$(5) \quad (\forall X \in \{0, 1\}^n) (\forall Y \in V) h(X, Y) = \prod_{A \subseteq N} \bigcup_{j=1}^n (z_{j,A,Y} + x_j),$$

where  $V = \{Y | Y \in \{0, 1\}^m \quad \wedge \quad b_\emptyset \prod_{\emptyset \neq S \subseteq N} (b_s + 1) = 0\}$ .

**Proof.** If  $h(X, Y) = \sum_{S \subseteq N} b_S \prod_{i \in S} x_i$  and  $b_\emptyset \prod_{\emptyset \neq S \subseteq N} (b_s + 1) = 0$ , it is obvious that a general solution

$$x_j = \sum_D g_{j,D} T^D \quad (j = 1, \dots, n)$$

of Boolean ring equation  $h(X, Y) = 0$  is of the form

$$x_j = \sum_D \left( \sum_C q_{j,D,C} Y^C \right) T^D \quad (j = 1, \dots, n),$$

since the coefficients  $g_{j,D}$  depend on  $Y$ , i.e.

$$g_{j,D} = \sum_C q_{j,D,C} Y^C \quad (j \in \{1, \dots, n\}, D \in \{0, 1\}^n)$$

by Theorem 3 and Theorem 1.

Taking in mind Theorem 7, we have the following equivalences:

$$\begin{aligned} & (\forall Y \in B^m) (b_{\emptyset} \prod_{\emptyset \neq S \subseteq N} (b_s + 1) = 0 \Rightarrow \\ & \quad (\forall X \in B^n) h(X, Y) = \prod_A \bigcup_{j=1}^n (\sum_D (\sum_C z_{j,D,C} Y^C) A^D + x_j) ) \\ \Leftrightarrow & (\forall Y \in B^m) (b_{\emptyset} \prod_{\emptyset \neq S \subseteq N} (b_s + 1) = 0 \Rightarrow \\ & \quad (\forall X \in B^n) h(X, Y) + \prod_A \bigcup_{j=1}^n (\sum_D (\sum_C z_{j,D,C} Y^C) A^D + x_j) = 0 ) \end{aligned}$$

(because  $(\forall a, b \in B) (a = b \Leftrightarrow a + b = 0)$ )

$$\begin{aligned} \Leftrightarrow & (\forall X \in B^m) (\forall Y \in B^n) (b_{\emptyset} \prod_{\emptyset \neq S \subseteq N} (b_s + 1) = 0 \Rightarrow \\ & \quad h(X, Y) + \prod_A \bigcup_{j=1}^n (\sum_D (\sum_C z_{j,D,C} Y^C) A^D + x_j) = 0 ) \\ \Leftrightarrow & (\forall X \in \{0, 1\}^n) (\forall Y \in \{0, 1\}^m) (b_{\emptyset} \prod_{\emptyset \neq S \subseteq N} (b_s + 1) = 0 \Rightarrow \\ & \quad h(X, Y) + \prod_A \bigcup_{j=1}^n (\sum_D (\sum_C z_{j,D,C} Y^C) A^D + x_j) = 0 ) \end{aligned}$$

(by Theorem 3 and Lemma)

$$\begin{aligned} \Leftrightarrow & (\forall X \in \{0, 1\}^n) (\forall Y \in V) h(X, Y) + \prod_A \bigcup_{j=1}^n (\sum_D (\sum_C z_{j,D,C} (Y) A^D + x_j) = 0 \\ \Leftrightarrow & (\forall X \in \{0, 1\}^n) (\forall Y \in V) h(X, Y) + \prod_A \bigcup_{j=1}^n (\sum_D (\sum_C z_{j,D,C} Y^C) A^D + x_j) \\ \Leftrightarrow & (\forall X \in \{0, 1\}^n) (\forall Y \in V) h(X, Y) = \prod_A \bigcup_{j=1}^n (\sum_C z_{j,A,C} Y^C + x_j) ) \\ \Leftrightarrow & (\forall X \in \{0, 1\}^n) (\forall Y \in V) h(X, Y) = \prod_A \bigcup_{j=1}^n (z_{j,A,Y} + x_j). \end{aligned}$$

If we solve the system (5) we get the general solution (4).

If we take  $Y^* \in V$ , we get the system of  $2^n$  equations:

$$(6) \quad (\forall X \in \{0, 1\}^n) h(X, Y^*) = \prod_A \bigcup_{j=1}^n (z_{j,A,Y^*} + x_j).$$

**Remark 1.**

(a) The system (6) does not contain the unknowns occurring in other equations of the system (5).

(b) Let  $S_h$  be the solution set of  $h(X, Y^*) = 0$ . If we take

$$\{(z_{1,A,Y^*}, \dots, z_{n,A,Y^*}) | A \in \{0, 1\}^n\} = S_h,$$

then the equation

$$h(X, Y^*) = \prod_A \bigcup_{j=1}^n (z_{j,A,Y^*} + x_j)$$

is satisfied. Namely, if  $h(X, Y^*) = 0$ , then

$$(z_{1,A,Y^*}, \dots, z_{n,A,Y^*}) = (x_1, \dots, x_n) \text{ for some } A \in \{0, 1\}^n,$$

i.e.

$$z_{j,A,Y^*} = x_j \quad (j = 1, \dots, n) \text{ for some } A \in \{0, 1\}^n,$$

i.e.

$$\bigcup_{j=1}^n (z_{j,A,Y^*} + x_j^*) = 0 \text{ for some } A \in \{0, 1\}^n,$$

i.e.

$$\prod_A \bigcup_{j=1}^n (z_{j,A,Y^*} + x_j^*) = 0.$$

(c) If  $Y_0 \notin V$ , then  $(z_{1,A,Y_0}, \dots, z_{n,A,Y_0})$  ( $A \in \{0, 1\}^n$ ) can be arbitrary element from  $\{0, 1\}^n$  because the  $n$ -tuple  $(z_{1,A,Y_0}, \dots, z_{n,A,Y_0})$  does not occur in (6).

The previous Remark 1 gives simple algorithm for solving the system (5).

If we use the known methods for solving Boolean ring equations ([9]), we really solve these equations in a Boolean ring  $\mathcal{B}$  or in some ring  $\mathcal{B}'$  generated by the coefficient appearing in these equations ([7]). Our Theorem 10 reduces the finding of a general solution of Boolean ring equation to solving simple equations in two element Boolean ring, i.e. in  $\{0, 1\}$ .

**Example .** Determine a general solution of the equation

$$axy + ay + b = 0$$

in arbitrary Boolean ring with unit.

Note that  $V = \{(0, 0), (1, 0), (1, 1)\}$  because of Theorem 9. Let

$$\begin{aligned} g_1(t_1, t_2, a, b) &= (p_{0,0}a'b' + p_{0,1}a'b + p_{0,2}ab' + p_{0,3}ab)t_1't_2' \\ &= (p_{1,0}a'b' + p_{1,1}a'b + p_{1,2}ab' + p_{1,3}ab)t_1't_2' \\ &= (p_{2,0}a'b' + p_{2,1}a'b + p_{2,2}ab' + p_{2,3}ab)t_1't_2' \\ &= (p_{3,0}a'b' + p_{3,1}a'b + p_{3,2}ab' + p_{3,3}ab)t_1't_2' \\ g_2(t_1, t_2, a, b) &= (q_{0,0}a'b' + q_{0,1}a'b + q_{0,2}ab' + q_{0,3}ab)t_1't_2' \\ &= (q_{1,0}a'b' + q_{1,1}a'b + q_{1,2}ab' + q_{1,3}ab)t_1't_2' \\ &= (q_{2,0}a'b' + q_{2,1}a'b + q_{2,2}ab' + q_{2,3}ab)t_1't_2' \\ &= (q_{3,0}a'b' + q_{3,1}a'b + q_{3,2}ab' + q_{3,3}ab)t_1't_2'. \end{aligned}$$



The system (5) becomes

$$axy + ay + b = \prod_{i=0}^3 ((p_{i,r} + x) \cup (q_{i,r} + y)) \quad (r \in \{0, 2, 3\}, (x, y) \in \{0, 1\}^2),$$

i.e.

$$\begin{aligned} 0 &= \prod_{i=0}^3 ((p_{i,0} + x) \cup (q_{i,0} + y)) \quad (x, y) \in \{0, 1\}^2 \\ xy + y &= \prod_{i=0}^3 ((p_{i,2} + x) \cup (q_{i,2} + y)) \quad (x, y) \in \{0, 1\}^2 \\ xy + y + 1 &= \prod_{i=0}^3 ((p_{i,3} + x) \cup (q_{i,3} + y)) \quad (x, y) \in \{0, 1\}^2. \end{aligned}$$

Let us introduce the notations

$$\begin{aligned} R_0 &= \{(p_{0,0}, q_{0,0}), (p_{1,0}, q_{1,0}), (p_{2,0}, q_{2,0}), (p_{3,0}, q_{3,0})\} \\ R_1 &= \{(p_{0,1}, q_{0,1}), (p_{1,1}, q_{1,1}), (p_{2,1}, q_{2,1}), (p_{3,1}, q_{3,1})\} \\ R_2 &= \{(p_{0,2}, q_{0,2}), (p_{1,2}, q_{1,2}), (p_{2,2}, q_{2,2}), (p_{3,2}, q_{3,2})\} \\ R_3 &= \{(p_{0,3}, q_{0,3}), (p_{1,3}, q_{1,3}), (p_{2,3}, q_{2,3}), (p_{3,3}, q_{3,3})\} \end{aligned}$$

In accordance with Remark 1 (c),  $R_1$  contains arbitrary elements from the set  $\{0, 1\}^2$ . Further, the solutions sets of the equations  $0 = 0$ ,  $xy + y = 0$  and  $xy + y + 1 = 0$  are  $\{(0, 0), (0, 1), (1, 0), (1, 1)\}$ ,  $\{(0, 0), (1, 0), (1, 1)\}$  and  $\{(0, 1), (1, 0)\}$ , respectively. Having in mind Remark 1 (b) we can take, for instance,

$$\begin{aligned} \{(p_{0,0}, q_{0,0}), (p_{1,0}, q_{1,0}), (p_{2,0}, q_{2,0}), (p_{3,0}, q_{3,0})\} &= \{(0, 0), (0, 1), (1, 0), (1, 1)\} \\ R_1 &= \{(p_{0,1}, q_{0,1}), (p_{1,1}, q_{1,1}), (p_{2,1}, q_{2,1}), (p_{3,1}, q_{3,1})\} = \{(0, 0), (0, 0), (0, 0), (0, 0)\} \\ R_2 &= \{(p_{0,2}, q_{0,2}), (p_{1,2}, q_{1,2}), (p_{2,2}, q_{2,2}), (p_{3,2}, q_{3,2})\} = \{(0, 0), (0, 0), (1, 0), (1, 1)\} \\ R_3 &= \{(p_{0,3}, q_{0,3}), (p_{1,3}, q_{1,3}), (p_{2,3}, q_{2,3}), (p_{3,3}, q_{3,3})\} = \{(0, 1), (0, 1), (1, 0), (0, 1)\} \end{aligned}$$

Thus, a general solution is determined by

$$\begin{aligned} x &= (a'b' + ab')t'_1 t_2 + (a'b' + ab')t_1 t_2 \\ y &= abt'_1 t_2 + (a'b' + ab)t'_1 t_2 + abt_1 t'_2 + (a'b' + ab' + ab)t_1 t_2, \end{aligned}$$

i.e.

$$\begin{aligned} x &= (b + 1)t_1(t_2 + 1) + (b + 1)t_1 t_2 \\ y &= ab(t_1 + 1)(t_2 + 1) + (a + b + 1)(t_1 + 1)t_2 + abt_1(t_2 + 1) + (ab + b + 1)t_1 t_2. \end{aligned}$$

# References

- [1] D. Banković. Some remarks on the number of the parameters of the solutions of Boolean equations. *Discrete Mathematics*, **79**, 1989/90, 229-234.
- [2] D. Banković. A note on Boolean equations. *Bulletin de la Societe Math. de Belgique*, **41**, 1989, 1, ser B, 47-53.
- [3] D. Banković. Certain Boolean equation. *Discrete Applied Mathematics*, **35**, 1992, 21-27.
- [4] J. P. Seschamps. Parametric solutions of Boolean equations. *Discrete Mathematics*, **3**, 1972, 333-342.
- [5] C. Ghilezan. Une généralisation du théorème de Löwenheim sur les équations de Boole. *Publ. Inst. Math. Beograd*, **11 (25)**, 1971, 57-59.
- [6] L. Löwenheim. Über die Auflösung von Gleichungen im logischen Gebietkalkul. *Sitzungber. Berl. Math. Gesel.*, **7**, 1910, 89-94.
- [7] U. Martin, T. Nipkow. Unification in Boolean rings. *Journal of Automated Reasoning*, **4**, 1988, 381-396.
- [8] S.B. Prešić. Ein Satz über reproductive Lösungen. *Publ. Inst. Math. Beograd*, **14(28)**, 1973, 133-136.
- [9] S. Rudeanu. *Boolean Functions and Equations*, North-Holland, 1974.

Faculty of Science  
 Radoja Domanovića 12  
 34 000 Kragujevac  
 YUGOSLAVIA

Received 29.12.1992