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# New Extremal Self-Dual Codes of Length 66

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Presented by P. Kenderov

A number of new extremal self-dual codes of length 66 having weight enumerators for which extremal self-dual codes were not previously known to exist, are constructed.

## 1. Introduction

A (binary) [n,k] code C is a k-dimensional vector subspace of  $\mathbb{F}_2^n$ , where  $\mathbb{F}_2$  is the field of two elements. An [n,k,d] code is an [n,k] code with minimum weight d. A code C is self-dual if  $C = C^{\perp}$  where  $C^{\perp}$  is the dual code of C. A self-dual code C is doubly-even if all codewords of C have weight divisible by four, and singly-even if there is at least one codeword of weight  $\equiv 2 \pmod{4}$ . Note that a doubly-even self-dual code of length n exists if and only if n is divisible by eight. An automorphism of C is a permutation of the coordinates of C which preserves C. The set consisting of all automorphisms of C forms a group called the automorphism group of C.

A singly-even self-dual code is called extremal if it has the largest minimum weight among all singly-even self-dual codes of that length. Conway and Sloane [1] proved a new upper bound for the minimum weight of singly-even self-dual codes and gave a list of the possible weight enumerators of extremal singly-even self-dual codes for lengths  $n \leq 64$  and length 72. Their work was extended to lengths up to 100 in [3]. Much work has been done concerning construction of extremal singly-even self-dual codes (see e.g., [1], [2], [3], [4], [8], [9], [10], [11], [14] and the references given therein). The largest minimum weight among self-dual code of length 66 is 12. The possible weight enumerators of

extremal self-dual [66, 33, 12] codes are

$$W_1 = 1 + (858 + 8\alpha)y^{12} + (18678 - 24\alpha)y^{14} + \cdots \quad (0 \le \alpha \le 778)$$

$$W_2 = 1 + 1690y^{12} + 7990y^{14} + \cdots$$

$$W_3 = 1 + (858 + 8\alpha)y^{12} + (18166 - 24\alpha)y^{14} + \cdots \quad (14 \le \alpha \le 756),$$

where  $\alpha$  is a parameter [3]. An extremal self-dual code D16 in [1] has weight enumerator  $W_1$  with  $\alpha = 0$ . It is shown in [4] that there are exactly three inequivalent double circulant self-dual [66, 33, 12] codes. Two of these have weight enumerator  $W_1$  with  $\alpha = 0$ , and the other has weight enumerator  $W_1$  with  $\alpha = 66$ . Tsai [12] constructed two inequivalent extremal self-dual codes with weight enumerator  $W_2$ . The construction of extremal self-dual codes of length 66 is lacked comparing with the construction of the codes of lengths  $\leq 64$ .

In this note, we construct a number of extremal self-dual [66, 33, 12] codes with weight enumerators  $W_1$  where  $\alpha = 3, 8, 10, 14, 15, 16, 17, 22, 24, 26, 31, 36, 38, 41, 43, 45, 46, 52, 59, 73, 74, 76, 78, 80, by several construction methods. These codes have weight enumerators for which extremal self-dual codes were not previously known to exist. A number of new extremal self-dual codes with weight enumerators <math>W_2$  are also constructed. The results in this note were already announced in the recent survey on self-dual codes written by Huffman [7].

# 2. Construction I

Let  $G_0 = (I_n, A)$  be a generator matrix of a self-dual code of length 2n, where  $I_n$  denotes the identity matrix of order n. Let  $(x_1, x_2, \ldots, x_n)$  be a vector of length n and odd (resp. even) weight if n is even (resp. odd). Then the following matrix

$$G = \begin{pmatrix} 1 & 0 & x_1 & \cdots & x_n & 1 & \cdots & 1 \\ y_1 & y_1 & & & & & \\ \vdots & \vdots & & I_n & & & A \\ y_n & y_n & & & & & \end{pmatrix}$$

where  $y_i = x_i + 1$   $(1 \le i \le n)$ , generates a self-dual code C of length 2n + 2 [5]. In this section, we construct new extremal self-dual codes of length 66 using this construction method.

A code having a generator matrix of the form ( $I_n$ , R) is called (pure) double circulant, where R is an  $n \times n$  circulant matrix. It was shown in [4] that there exist exactly two extremal double circulant singly-even self-dual [64, 32, 12] codes, up to equivalence. These codes have generator matrices  $M_1$  and  $M_2$  with

the following first rows for R

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(0001011111011100011111111111100011) and (00000001000101001010010111111000111),
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respectively. Recently an extremal singly-even self-dual [64, 32, 12] code has been found in [10]. A generator matrix of the form (I, A) of the code is given in [10] and we denote the matrix A by  $M_3$ .

By the above method, from  $M_1, M_2$  and  $M_3$ , we have found new extremal self-dual [66, 33, 12] codes having weight enumerators for which extremal self-dual codes were not previously known to exist, where the results are listed in Table 1. In the table, we list the used generator matrices  $G_0$ , the vectors  $x = (x_1, x_2, \ldots, x_{32})$ , the weight enumerators W and the orders |Aut| of the automorphism groups for the new codes. The automorphism groups are calculated by MAGMA.

Codes	$G_0$	Vectors $x$	W	Aut
$\overline{C_1}$	$M_1$	(11111010001000011000000010111000)	$W_1 \ (\alpha = 41)$	1
$C_2$	$M_1$	(11100010001000000000101001100000)	$W_1 \ (\alpha = 43)$	1
$C_3$	$M_1$	(11110001000000100000001100010110)	$W_1 \ (\alpha = 46)$	2
$\overline{C_4}$	$M_2$	(1101011001011000000000000000010000)	$W_1 \ (\alpha = 74)$	2
$C_5$	$M_2$	(11000011000010000010101000100000)	$W_1 \ (\alpha = 76)$	1
$C_6$	$M_2$	(111001001000010001000000000010100)	$W_1 \ (\alpha = 78)$	1
$\overline{C_7}$	$M_3$	(10000010110000101010000100100101)	$W_1 \ (\alpha = 8)$	1
$C_8$	$M_3$	(10000101100010010100001010001100)	$W_1 \ (\alpha = 14)$	1
$C_9$	$M_3$	(00010000101001100100010001000010)	$W_1 \; (\alpha = 16)$	1
$C_{10}$	$M_3$	(0000100010010010010011000110000)	$W_1 \; (\alpha = 26)$	1

Table 1: New extremal self-dual [66, 33, 12] codes I

#### 3. Construction II

In this section, we investigate extremal self-dual [66, 33, 12] codes with an automorphism of order 7 using the technique developed by Huffman [6] and Yorgov [13]. Extremal singly-even self-dual codes of lengths 42, 50, 52, 54 and 60, 64 with automorphisms of order 7 have been investigated in [14], [8], [9], [11] and [2], respectively. In this section, we extend the work to length 66.

Suppose that  $\sigma$  is an automorphism of order 7 of an extremal self-dual [66, 33, 12] code D. By [13, Theorem 1], one can show that  $\sigma$  consists of 9 7-cycles together with 3 fixed points. Let  $E_{\sigma}(D)$  be the set of the vectors in D, which have even weight in each cycle of  $\sigma$ . Set  $F_{\sigma}(D) = \{v \in D \mid v\sigma = v\}$ . By

[6, Theorem 1], it follows  $D = F_{\sigma}(D) \oplus E_{\sigma}(D)$  where  $\oplus$  denotes the direct sum. Each vector v from  $F_{\sigma}(D)$  is constant on any cycle of  $\sigma$ . Using the results from [2] we construct a possible generator matrix of the code D of the form:

(1) 
$$G_{66} = \begin{pmatrix} X & \\ Z & O \end{pmatrix},$$

where X is a  $6 \times 66$  matrix generating the fixed subcode  $F_{\sigma}(D)$  under  $\sigma$  of D and O is the  $27 \times 3$  zero-matrix. Note that Z has the following form:

where the cells  $r_i$ ,  $k_i$ , for i = 0, ..., 20, are  $3 \times 7$  (right) circulant matrices with first rows corresponding to polynomials in  $\mathbb{F}_2[x]/(x^7-1)$  and where the blanks are equal to the zero.

We have found many matrices of the form (1) which generate extremal self-dual [66, 33, 12] codes. Here we present 11 codes  $D_i$ , (i = 1, 2, ..., 11), with weight enumerators for which extremal self-dual codes were not previously known to exist. For each code  $D_i$ , we denote the matrices X and Z in the generator matrices (1) by  $X_i$  and  $Z_i$ , respectively. In Tables 2 and 3, we give the matrices  $X_i$  and the cells of  $Z_i$ , respectively. In Table 2, a is the all-one vector of length 7 and the blanks are zero's, and in Table 3  $u_i$  for i = 0, ..., 6 and  $v_i$  for i = 0, ..., 6, are  $3 \times 7$  (right) circulant matrices with first rows 1110100, 0111010, 0011101, 10011110, 0100111, 1010011, 1100011, 1100101, 10111001, 01011110 and 0010111 respectively, whereas the cell 0 denotes the  $3 \times 7$  zero-matrix.

For the 11 codes  $D_i$ , we list the weight enumerators W and the orders |Aut| of the automorphism groups in Table 4. Note that an extremal self-dual code with weight enumerator  $W_1$  ( $\alpha = 66$ ) is also found. Since the double circulant code with the same weight enumerator in [4] has no automorphism of order 7, the codes are inequivalent.

	a	a		a						U	1	U	ll l	a	a	a	a						U	U	U
	a	a	a							1	0	0				a	a	a					1	0	0
3.5			a		a	a				1	0	0						a	a	a			1	0	0
$X_1$					a	a	a			0	0	1	$X_7$						a	a	a		0	1	0
							a	a	a	0	0	1									a	a	0	1	1
		a				a	_	-	a	1	1	1			a		a			a	_	a	1	1	0
	a	a	a							1	0	0		a	a	a	a						0	0	0
	ω.		a	a	a					1	ő	ő		a		a	a	a					1	ő	ő
				a	a	a				0	1	ő				· ·		a	a	a			1	ő	ő
$X_2$					a	a	a			ő	1	1	$X_8$						a	a	a		ō	1	ő
							a	a	a	ő	0	1									a	a	ŏ	1	1
		a			a		a	-	a	1	1	0			a		a			a	_	a	1	1	0
	a	a			-					1	1	0		a	a	a							1	0	0
	_	a	a	a						0	1	0			a	a	a						0	1	Õ
			a	a	a	a				Õ	0	0					a	a					Õ	1	1
$X_3$					a	a	a	a		0	0	0	$X_9$					a	a	a			0	0	1
							a	a	a	Õ	Õ	1							a	a	a	a	Õ	ō	0
				a		a		a	a	1	1	0				a				a		a	1	1	1
	a	a	a							1	0	0		a	a	a							1	0	0
		a	a	a						0	1	0			a	a	a						0	1	0
3.5				a	a	a				0	1	0					a	a					0	1	1
$X_4$					a	a	a			0	0	1	$X_{10}$					a	a	a			0	0	1
							a	a	a	0	0	1							a	a	a	a	0	0	0
			a			a			a	1	1	1				a		a		a		a	1	1	0
	a	a								1	1	0			a	a	a	a					0	0	0
		a	a	a						0	1	0		a		a	a						1	0	0
3.5			a	a	a					0	0	1		a					a				1	1	0
$X_5$					a	a	a			0	0	1	$X_{11}$						a	a			0	1	1
						a	a	a	a	0	0	0								a	a	a	0	0	1
				a			a		a	1	1	1		a	a		a			a		a	0	1	0
	a	a	a	a						0	0	0													
			a	a	a	a				0	0	0													
1/					a	a	a			1	0	0													
$X_6$							a	a		1	1	0													
								a	a	0	1	1													
		a		a		a				1	1	1													

Table 2: Matrices  $X_i$ 

### 4. Construction III

Two self-dual codes C and C' of length n are called *neighbors* if the dimension of  $C \cap C'$  is n/2 - 1. Let  $v \in \mathbb{F}_2^{66}$  be a vector of even weight and let C be a self-dual code of length 66. For v and C ( $v \notin C$ ), we define N(v,C) as follows:

$$N(v,C) = (C \cap \langle v \rangle^{\perp}) \cup \{u + v \mid u \in (C \cap \langle v \rangle^{\perp})\}.$$

Then it is well known that N(v,C) is a self-dual code which is a neighbor of C. As described in Section 1,the code D16 in [1] is an extremal self-dual double circulant code with generator matrix ( $I_{33}$ , R) where R is the circulant matrix with first row

## (0000010110010110110010111111011001).

Using the above neighbor construction, we have found new extremal self-dual codes having weight enumerators for which extremal self-dual codes were not previously known to exist. For the new self-dual neighbors N(C, v), C, v, the weight enumerators W and the orders |Aut| of the automorphism groups are

Table 3: Cells of the matrices  $Z_i$ 

Codes	$r_0, r_1, \ldots, r_{20}, k_0, k_1, \ldots, k_{20}$
$D_1$	$u_0, u_0, u_0, u_0, u_0, u_0, u_1, u_1, u_1, u_0, u_1, u_1, u_2, u_0, u_1, u_2, u_2, u_0, u_5, u_2, u_4$
	$v_0, v_0, v_0, v_0, v_0, v_0, v_0, v_1, v_1, v_1, v_5, v_0, v_1, v_1, v_2, v_2, v_0, v_1, v_2, v_2, v_4$
$D_2$	$u_0, u_0, u_0, u_0, u_0, u_0, 0, 0, u_0, u_$
	$v_0, v_0, v_0, v_0, v_0, v_0, v_0, v_0, $
$D_3$	$u_0, u_0, u_0, u_0, u_0, u_0, 0, 0, u_0, u_$
	$v_0, v_0, v_0, v_0, v_0, v_0, v_0, 0, 0, 0, v_0, v_$
$D_4$	$u_0, u_0, u_0, u_0, u_0, u_0, 0, 0, u_0, u_$
	$v_0, v_0, v_0, v_0, v_0, v_0, v_0, 0, 0, 0, v_2, v_0, 0, 0, v_0, v_3, v_0, v_0, v_1, v_2, v_2$
$D_5$	$u_0, u_0, u_0, u_0, u_0, u_0, 0, 0, u_0, u_$
	$v_0, v_0, v_0, v_0, v_0, v_0, v_0, 0, 0, 0, v_0, v_$
$D_6$	$u_0, u_0, u_0, u_0, u_0, u_0, 0, 0, u_0, u_$
	$v_0, v_0, v_0, v_0, v_0, v_0, v_0, 0, 0, 0, v_1, v_0, 0, 0, v_0, v_2, v_0, v_0, v_1, v_2, v_3$
$D_7$	$u_0, u_0, u_0, u_0, u_0, u_0, 0, 0, u_0, u_$
	$v_0, v_0, v_0, v_0, v_0, v_0, v_0, 0, 0, 0, v_1, v_0, 0, 0, v_0, v_5, v_0, v_0, v_1, v_2, v_3$
$D_8$	$u_0, u_0, u_0, u_0, u_0, u_0, 0, 0, u_0, u_$
	$v_0, v_0, v_0, v_0, v_0, v_0, v_0, 0, 0, 0, v_2, v_0, 0, 0, v_0, v_5, v_0, v_0, v_1, v_2, 0$
$D_9$	$u_0, u_0, u_0, u_0, u_0, u_0, 0, 0, u_0, u_$
	$v_0, v_0, v_0, v_0, v_0, v_0, v_0, 0, 0, 0, v_3, v_0, 0, 0, v_0, v_2, v_0, v_0, v_1, v_2, v_6$
$D_{10}$	$u_0, u_0, u_0, u_0, u_0, u_0, 0, 0, u_0, u_$
	$v_0, v_0, v_0, v_0, v_0, v_0, v_0, 0, 0, 0, v_1, v_0, 0, 0, v_0, 0, v_0, v_0, v_1, v_4, v_2$
$D_{11}$	$u_0, u_0, u_0, u_0, u_0, u_0, u_1, u_1, u_1, u_0, u_1, u_1, u_2, u_0, u_1, u_2, u_3, u_0, u_1, u_2, u_5$
	$v_0, v_0, v_0, v_0, v_0, v_0, v_0, v_1, v_1, v_1, v_1, v_0, v_1, v_1, v_2, v_2, v_0, v_1, v_2, v_3, v_5$

listed in Table 5 where

 $v_2 = (1001000001000000010000100000100$ 

011011100010110111000101101110001).

Table 4: New extremal self-dual [66, 33, 12] codes II

Codes	W	Aut	Codes	W	Aut
$D_1$	$W_1 \ (\alpha = 3)$	7	$D_2$	$W_1 \ (\alpha = 10)$	7
$D_3$	$W_1 \ (\alpha = 17)$	7	$D_4$	$W_1 \ (\alpha = 24)$	7
$D_5$	$W_1 \ (\alpha = 31)$	7	$D_6$	$W_1 \ (\alpha = 38)$	7
$D_7$	$W_1 \ (\alpha = 45)$	7	$D_8$	$W_1 \ (\alpha = 52)$	7
$D_9$	$W_1 \ (\alpha = 59)$	7	$D_{10}$	$W_1 \ (\alpha = 73)$	7
$D_{11}$	$W_1 \ (\alpha = 80)$	7			

Table 5: New extremal self-dual [66, 33, 12] codes III

Codes	C	v	W	Aut
$E_1$	D16	$v_1$	$W_1 \ (\alpha = 15)$	15
$E_2$	$C_{66,25}$ in [4]	$v_2$	$W_1 \ (\alpha = 22)$	220
$E_3$	$C_{66,25}$ in [4]	$v_3$	$W_1 \ (\alpha = 36)$	60

As described in Section , Tsai [12] constructed two codes with the weight enumerator  $W_2$ . The two codes are denoted by T66 and T66' in [12]. We remark that both T66 and T66' are neighbors of D16. In fact, if

v = (110101101110001000111011000001111

11110101011111111100011010010010011)

then we find T66 = N(D16, v).

Table 6: Extremal self-dual [66, 33, 12] codes with weight enumerator  $W_2$ 

41024713577, 41024737517, 41034175676, 41034512156, 41067152331, 41067174071 $41067573512,\,41076335057,\,41124157354,\,41125151432,\,41125556333,\,45476332745$  $45663336673,\ 45732732567,\ 46427334742,\ 46433352450,\ 46433754565,\ 46436310766$ 46476332475, 47473730744, 47522336751, 47536312452, 47736710561, 47737731462 54426371640, 54476713763, 61136775505, 61177060726, 61327061524, 61377243367 $61537742727,\, 61727342776,\, 61736554744,\, 61767444306,\, 63126252505,\, 63126741144$  $63166570374,\ 63176256337,\ 63177564356,\ 63536172756,\ 63566661714,\ 63566670155$  $63567652115,\ 63726274537,\ 63776467126,\ 65126377115,\ 65126574146,\ 65136142154$  $65137042327,\ 65137142127,\ 65336067714,\ 65337744347,\ 65366144507,\ 65377554127$  $65532332757,\ 65536546564,\ 65572374550,\ 65732332455,\ 65732716644,\ 65736354544$  $65736455704,\ 65766544744,\ 66562730655,\ 67136771526,\ 67166254146,\ 67167574316$  $67177342137,\ 67526357441,\ 71137264374,\ 71166364107,\ 71326250556,\ 71337344357$  $71526752567,\ 71527353566,\ 71537541727,\ 73127741727,\ 73176466726,\ 73326052575$  $75177041104,\ 75326367516,\ 75366045104,\ 75567061106,\ 75727253147,\ 76532713454$ 77136075347, 77166052745, 77366471575, 77377066126, 77526333466, 7752675250777563375474, 77726166156, 77766045505, 77767156736

Moreover, we have found a number of new extremal self-dual codes N(T66, v) with the weight enumerator  $W_2$  as self-dual neighbors of T66 under the assumption that v has the form  $(v_1, v_2, \ldots, v_{33}, 0, \ldots, 0)$ . In Table 6, we give the vectors  $(v_1, v_2, \ldots, v_{33})$  for 100 self-dual neighbors where the vectors  $(v_1, v_2, \ldots, v_{33})$  are written in octal using  $0 = (000), 1 = (001), \ldots, 6 = (110)$ 

and 7 = (111) to save space. Using MAGMA, we have verified that the 100 codes are inequivalent and have a trivial automorphism group. Since we have verified that the codes T66 and T66' have an automorphism group of order 5, our codes are inequivalent to the two codes.

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