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New Extremal Self-Dual Codes of Length 66

Masaaki Harada ¹, Takuji Nishimura ¹ and Radinka Yorgova ²

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A number of new extremal self-dual codes of length 66 having weight enumerators for which extremal self-dual codes were not previously known to exist, are constructed.

1. Introduction

A (binary) [n,k] code C is a k-dimensional vector subspace of \mathbb{F}_2^n , where \mathbb{F}_2 is the field of two elements. An [n,k,d] code is an [n,k] code with minimum weight d. A code C is self-dual if $C = C^{\perp}$ where C^{\perp} is the dual code of C. A self-dual code C is doubly-even if all codewords of C have weight divisible by four, and singly-even if there is at least one codeword of weight $\equiv 2 \pmod{4}$. Note that a doubly-even self-dual code of length n exists if and only if n is divisible by eight. An automorphism of C is a permutation of the coordinates of C which preserves C. The set consisting of all automorphisms of C forms a group called the automorphism group of C.

A singly-even self-dual code is called extremal if it has the largest minimum weight among all singly-even self-dual codes of that length. Conway and Sloane [1] proved a new upper bound for the minimum weight of singly-even self-dual codes and gave a list of the possible weight enumerators of extremal singly-even self-dual codes for lengths $n \leq 64$ and length 72. Their work was extended to lengths up to 100 in [3]. Much work has been done concerning construction of extremal singly-even self-dual codes (see e.g., [1], [2], [3], [4], [8], [9], [10], [11], [14] and the references given therein). The largest minimum weight among self-dual code of length 66 is 12. The possible weight enumerators of

extremal self-dual [66, 33, 12] codes are

$$W_1 = 1 + (858 + 8\alpha)y^{12} + (18678 - 24\alpha)y^{14} + \cdots \quad (0 \le \alpha \le 778)$$

$$W_2 = 1 + 1690y^{12} + 7990y^{14} + \cdots$$

$$W_3 = 1 + (858 + 8\alpha)y^{12} + (18166 - 24\alpha)y^{14} + \cdots \quad (14 \le \alpha \le 756),$$

where α is a parameter [3]. An extremal self-dual code D16 in [1] has weight enumerator W_1 with $\alpha = 0$. It is shown in [4] that there are exactly three inequivalent double circulant self-dual [66, 33, 12] codes. Two of these have weight enumerator W_1 with $\alpha = 0$, and the other has weight enumerator W_1 with $\alpha = 66$. Tsai [12] constructed two inequivalent extremal self-dual codes with weight enumerator W_2 . The construction of extremal self-dual codes of length 66 is lacked comparing with the construction of the codes of lengths ≤ 64 .

In this note, we construct a number of extremal self-dual [66, 33, 12] codes with weight enumerators W_1 where $\alpha = 3, 8, 10, 14, 15, 16, 17, 22, 24, 26, 31, 36, 38, 41, 43, 45, 46, 52, 59, 73, 74, 76, 78, 80, by several construction methods. These codes have weight enumerators for which extremal self-dual codes were not previously known to exist. A number of new extremal self-dual codes with weight enumerators <math>W_2$ are also constructed. The results in this note were already announced in the recent survey on self-dual codes written by Huffman [7].

2. Construction I

Let $G_0 = (I_n, A)$ be a generator matrix of a self-dual code of length 2n, where I_n denotes the identity matrix of order n. Let (x_1, x_2, \ldots, x_n) be a vector of length n and odd (resp. even) weight if n is even (resp. odd). Then the following matrix

$$G = \begin{pmatrix} 1 & 0 & x_1 & \cdots & x_n & 1 & \cdots & 1 \\ y_1 & y_1 & & & & & \\ \vdots & \vdots & & I_n & & & A \\ y_n & y_n & & & & & \end{pmatrix}$$

where $y_i = x_i + 1$ $(1 \le i \le n)$, generates a self-dual code C of length 2n + 2 [5]. In this section, we construct new extremal self-dual codes of length 66 using this construction method.

A code having a generator matrix of the form (I_n , R) is called (pure) double circulant, where R is an $n \times n$ circulant matrix. It was shown in [4] that there exist exactly two extremal double circulant singly-even self-dual [64, 32, 12] codes, up to equivalence. These codes have generator matrices M_1 and M_2 with

the following first rows for R

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(0001011111011100011111111111100011) and (00000001000101001010010111111000111),
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respectively. Recently an extremal singly-even self-dual [64, 32, 12] code has been found in [10]. A generator matrix of the form (I, A) of the code is given in [10] and we denote the matrix A by M_3 .

By the above method, from M_1, M_2 and M_3 , we have found new extremal self-dual [66, 33, 12] codes having weight enumerators for which extremal self-dual codes were not previously known to exist, where the results are listed in Table 1. In the table, we list the used generator matrices G_0 , the vectors $x = (x_1, x_2, \ldots, x_{32})$, the weight enumerators W and the orders |Aut| of the automorphism groups for the new codes. The automorphism groups are calculated by MAGMA.

Codes	G_0	Vectors x	W	Aut
$\overline{C_1}$	M_1	(11111010001000011000000010111000)	$W_1 \ (\alpha = 41)$	1
C_2	M_1	(11100010001000000000101001100000)	$W_1 \ (\alpha = 43)$	1
C_3	M_1	(11110001000000100000001100010110)	$W_1 \ (\alpha = 46)$	2
$\overline{C_4}$	M_2	(1101011001011000000000000000010000)	$W_1 \ (\alpha = 74)$	2
C_5	M_2	(11000011000010000010101000100000)	$W_1 \ (\alpha = 76)$	1
C_6	M_2	(111001001000010001000000000010100)	$W_1 \ (\alpha = 78)$	1
$\overline{C_7}$	M_3	(10000010110000101010000100100101)	$W_1 \ (\alpha = 8)$	1
C_8	M_3	(10000101100010010100001010001100)	$W_1 \ (\alpha = 14)$	1
C_9	M_3	(00010000101001100100010001000010)	$W_1 \; (\alpha = 16)$	1
C_{10}	M_3	(0000100010010010010011000110000)	$W_1 \; (\alpha = 26)$	1

Table 1: New extremal self-dual [66, 33, 12] codes I

3. Construction II

In this section, we investigate extremal self-dual [66, 33, 12] codes with an automorphism of order 7 using the technique developed by Huffman [6] and Yorgov [13]. Extremal singly-even self-dual codes of lengths 42, 50, 52, 54 and 60, 64 with automorphisms of order 7 have been investigated in [14], [8], [9], [11] and [2], respectively. In this section, we extend the work to length 66.

Suppose that σ is an automorphism of order 7 of an extremal self-dual [66, 33, 12] code D. By [13, Theorem 1], one can show that σ consists of 9 7-cycles together with 3 fixed points. Let $E_{\sigma}(D)$ be the set of the vectors in D, which have even weight in each cycle of σ . Set $F_{\sigma}(D) = \{v \in D \mid v\sigma = v\}$. By

[6, Theorem 1], it follows $D = F_{\sigma}(D) \oplus E_{\sigma}(D)$ where \oplus denotes the direct sum. Each vector v from $F_{\sigma}(D)$ is constant on any cycle of σ . Using the results from [2] we construct a possible generator matrix of the code D of the form:

(1)
$$G_{66} = \begin{pmatrix} X & \\ Z & O \end{pmatrix},$$

where X is a 6×66 matrix generating the fixed subcode $F_{\sigma}(D)$ under σ of D and O is the 27×3 zero-matrix. Note that Z has the following form:

where the cells r_i , k_i , for i = 0, ..., 20, are 3×7 (right) circulant matrices with first rows corresponding to polynomials in $\mathbb{F}_2[x]/(x^7-1)$ and where the blanks are equal to the zero.

We have found many matrices of the form (1) which generate extremal self-dual [66, 33, 12] codes. Here we present 11 codes D_i , (i = 1, 2, ..., 11), with weight enumerators for which extremal self-dual codes were not previously known to exist. For each code D_i , we denote the matrices X and Z in the generator matrices (1) by X_i and Z_i , respectively. In Tables 2 and 3, we give the matrices X_i and the cells of Z_i , respectively. In Table 2, a is the all-one vector of length 7 and the blanks are zero's, and in Table 3 u_i for i = 0, ..., 6 and v_i for i = 0, ..., 6, are 3×7 (right) circulant matrices with first rows 1110100, 0111010, 0011101, 10011110, 0100111, 1010011, 1100011, 1100101, 10111001, 01011110 and 0010111 respectively, whereas the cell 0 denotes the 3×7 zero-matrix.

For the 11 codes D_i , we list the weight enumerators W and the orders |Aut| of the automorphism groups in Table 4. Note that an extremal self-dual code with weight enumerator W_1 ($\alpha = 66$) is also found. Since the double circulant code with the same weight enumerator in [4] has no automorphism of order 7, the codes are inequivalent.

	a	a		a						U	1	U	ll l	a	a	a	a						U	U	U
	a	a	a							1	0	0				a	a	a					1	0	0
3.5			a		a	a				1	0	0						a	a	a			1	0	0
X_1					a	a	a			0	0	1	X_7						a	a	a		0	1	0
							a	a	a	0	0	1									a	a	0	1	1
		a				a	_	-	a	1	1	1			a		a			a	_	a	1	1	0
	a	a	a							1	0	0		a	a	a	a						0	0	0
	ω		a	a	a					1	ő	ő		a		a	a	a					1	ő	ő
				a	a	a				0	1	ő				a		a	a	a			1	ő	ő
X_2					a	a	a			ŏ	1	1	X_8						a	a	a		ō	1	ő
							a	a	a	ŏ	0	1									a	a	ŏ	1	1
		a			a		a	-	a	1	1	0			a		a			a	_	a	1	1	0
	a	a			-					1	1	0		a	a	a							1	0	0
	_	a	a	a						0	1	0			a	a	a						0	1	Õ
			a	a	a	a				Õ	0	0					a	a					Õ	1	1
X_3					a	a	a	a		0	0	0	X_9					a	a	a			0	0	1
							a	a	a	Õ	Õ	1							a	a	a	a	Õ	ō	0
				a		a		a	a	1	1	0				a				a		a	1	1	1
	a	a	a							1	0	0		a	a	a							1	0	0
		a	a	a						0	1	0			a	a	a						0	1	0
3.5				a	a	a				0	1	0					a	a					0	1	1
X_4					a	a	a			0	0	1	X_{10}					a	a	a			0	0	1
							a	a	a	0	0	1							a	a	a	a	0	0	0
			a			a			a	1	1	1				a		a		a		a	1	1	0
	a	a								1	1	0			a	a	a	a					0	0	0
		a	a	a						0	1	0		a		a	a						1	0	0
3.5			a	a	a					0	0	1		a					a				1	1	0
X_5					a	a	a			0	0	1	X_{11}						a	a			0	1	1
						a	a	a	a	0	0	0								a	a	a	0	0	1
				a			a		a	1	1	1		a	a		a			a		a	0	1	0
	a	a	a	a						0	0	0													
			a	a	a	a				0	0	0													
1/					a	a	a			1	0	0													
X_6							a	a		1	1	0													
								a	a	0	1	1													
		a		a		a				1	1	1													

Table 2: Matrices X_i

4. Construction III

Two self-dual codes C and C' of length n are called *neighbors* if the dimension of $C \cap C'$ is n/2 - 1. Let $v \in \mathbb{F}_2^{66}$ be a vector of even weight and let C be a self-dual code of length 66. For v and C ($v \notin C$), we define N(v,C) as follows:

$$N(v,C) = (C \cap \langle v \rangle^{\perp}) \cup \{u + v \mid u \in (C \cap \langle v \rangle^{\perp})\}.$$

Then it is well known that N(v,C) is a self-dual code which is a neighbor of C. As described in Section 1,the code D16 in [1] is an extremal self-dual double circulant code with generator matrix (I_{33} , R) where R is the circulant matrix with first row

(0000010110010110110010111111011001).

Using the above neighbor construction, we have found new extremal self-dual codes having weight enumerators for which extremal self-dual codes were not previously known to exist. For the new self-dual neighbors N(C, v), C, v, the weight enumerators W and the orders |Aut| of the automorphism groups are

Table 3: Cells of the matrices Z_i

Codes	$r_0, r_1, \ldots, r_{20}, k_0, k_1, \ldots, k_{20}$
D_1	$u_0, u_0, u_0, u_0, u_0, u_0, u_1, u_1, u_1, u_0, u_1, u_1, u_2, u_0, u_1, u_2, u_2, u_0, u_5, u_2, u_4$
	$v_0, v_0, v_0, v_0, v_0, v_0, v_0, v_1, v_1, v_1, v_5, v_0, v_1, v_1, v_2, v_2, v_0, v_1, v_2, v_2, v_4$
D_2	$u_0, u_0, u_0, u_0, u_0, u_0, 0, 0, u_0, u_$
	$v_0, v_0, v_0, v_0, v_0, v_0, v_0, v_0, $
D_3	$u_0, u_0, u_0, u_0, u_0, u_0, 0, 0, u_0, u_$
	$v_0, v_0, v_0, v_0, v_0, v_0, v_0, 0, 0, 0, v_0, v_$
D_4	$u_0, u_0, u_0, u_0, u_0, u_0, 0, 0, u_0, u_$
	$v_0, v_0, v_0, v_0, v_0, v_0, v_0, 0, 0, 0, v_2, v_0, 0, 0, v_0, v_3, v_0, v_0, v_1, v_2, v_2$
D_5	$u_0, u_0, u_0, u_0, u_0, u_0, 0, 0, u_0, u_$
	$v_0, v_0, v_0, v_0, v_0, v_0, v_0, 0, 0, 0, v_0, v_$
D_6	$u_0, u_0, u_0, u_0, u_0, u_0, 0, 0, u_0, u_$
	$v_0, v_0, v_0, v_0, v_0, v_0, v_0, v_0, $
D_7	$u_0, u_0, u_0, u_0, u_0, u_0, 0, 0, u_0, u_$
	$v_0, v_0, v_0, v_0, v_0, v_0, v_0, 0, 0, 0, v_1, v_0, 0, 0, v_0, v_5, v_0, v_0, v_1, v_2, v_3$
D_8	$u_0, u_0, u_0, u_0, u_0, u_0, 0, 0, u_0, u_$
-	$v_0, v_0, v_0, v_0, v_0, v_0, v_0, 0, 0, 0, v_2, v_0, 0, 0, v_0, v_5, v_0, v_0, v_1, v_2, 0$
D_9	$u_0, u_0, u_0, u_0, u_0, u_0, 0, 0, u_0, u_$
	$v_0, v_0, v_0, v_0, v_0, v_0, v_0, 0, 0, 0, v_3, v_0, 0, 0, v_0, v_2, v_0, v_0, v_1, v_2, v_6$
D_{10}	$u_0, u_0, u_0, u_0, u_0, u_0, 0, 0, u_0, u_$
	$v_0, v_0, v_0, v_0, v_0, v_0, v_0, 0, 0, 0, v_1, v_0, 0, 0, v_0, 0, v_0, v_0, v_1, v_4, v_2$
D_{11}	$u_0, u_0, u_0, u_0, u_0, u_0, u_1, u_1, u_1, u_0, u_1, u_1, u_2, u_0, u_1, u_2, u_3, u_0, u_1, u_2, u_5$
	$v_0, v_0, v_0, v_0, v_0, v_0, v_0, v_1, v_1, v_1, v_1, v_0, v_1, v_1, v_2, v_2, v_0, v_1, v_2, v_3, v_5$

listed in Table 5 where

 $v_2 = (1001000001000000010000100000100$

011011100010110111000101101110001).

Table 4: New extremal self-dual [66, 33, 12] codes II

Codes	W	Aut	Codes	W	Aut
D_1	$W_1 \ (\alpha = 3)$	7	D_2	$W_1 \ (\alpha = 10)$	7
D_3	$W_1 \ (\alpha = 17)$	7	D_4	$W_1 \ (\alpha = 24)$	7
D_5	$W_1 \ (\alpha = 31)$	7	D_6	$W_1 \ (\alpha = 38)$	7
D_7	$W_1 \ (\alpha = 45)$	7	D_8	$W_1 \ (\alpha = 52)$	7
D_9	$W_1 \ (\alpha = 59)$	7	D_{10}	$W_1 \ (\alpha = 73)$	7
D_{11}	$W_1 \ (\alpha = 80)$	7			

Table 5: New extremal self-dual [66, 33, 12] codes III

Codes	C	v	W	Aut
E_1	D16	v_1	$W_1 \ (\alpha = 15)$	15
E_2	$C_{66,25}$ in [4]	v_2	$W_1 \ (\alpha = 22)$	220
E_3	$C_{66,25}$ in [4]	v_3	$W_1 \ (\alpha = 36)$	60

As described in Section , Tsai [12] constructed two codes with the weight enumerator W_2 . The two codes are denoted by T66 and T66' in [12]. We remark that both T66 and T66' are neighbors of D16. In fact, if

v = (110101101110001000111011000001111

11110101011111111100011010010010011)

then we find T66 = N(D16, v).

Table 6: Extremal self-dual [66, 33, 12] codes with weight enumerator W_2

41024713577, 41024737517, 41034175676, 41034512156, 41067152331, 41067174071 $41067573512,\,41076335057,\,41124157354,\,41125151432,\,41125556333,\,45476332745$ $45663336673,\ 45732732567,\ 46427334742,\ 46433352450,\ 46433754565,\ 46436310766$ 46476332475, 47473730744, 47522336751, 47536312452, 47736710561, 47737731462 54426371640, 54476713763, 61136775505, 61177060726, 61327061524, 61377243367 $61537742727,\, 61727342776,\, 61736554744,\, 61767444306,\, 63126252505,\, 63126741144$ $63166570374,\ 63176256337,\ 63177564356,\ 63536172756,\ 63566661714,\ 63566670155$ $63567652115,\ 63726274537,\ 63776467126,\ 65126377115,\ 65126574146,\ 65136142154$ $65137042327,\ 65137142127,\ 65336067714,\ 65337744347,\ 65366144507,\ 65377554127$ $65532332757,\ 65536546564,\ 65572374550,\ 65732332455,\ 65732716644,\ 65736354544$ $65736455704,\ 65766544744,\ 66562730655,\ 67136771526,\ 67166254146,\ 67167574316$ $67177342137,\ 67526357441,\ 71137264374,\ 71166364107,\ 71326250556,\ 71337344357$ $71526752567,\ 71527353566,\ 71537541727,\ 73127741727,\ 73176466726,\ 73326052575$ $75177041104,\ 75326367516,\ 75366045104,\ 75567061106,\ 75727253147,\ 76532713454$ 77136075347, 77166052745, 77366471575, 77377066126, 77526333466, 7752675250777563375474, 77726166156, 77766045505, 77767156736

Moreover, we have found a number of new extremal self-dual codes N(T66, v) with the weight enumerator W_2 as self-dual neighbors of T66 under the assumption that v has the form $(v_1, v_2, \ldots, v_{33}, 0, \ldots, 0)$. In Table 6, we give the vectors $(v_1, v_2, \ldots, v_{33})$ for 100 self-dual neighbors where the vectors $(v_1, v_2, \ldots, v_{33})$ are written in octal using $0 = (000), 1 = (001), \ldots, 6 = (110)$

and 7 = (111) to save space. Using MAGMA, we have verified that the 100 codes are inequivalent and have a trivial automorphism group. Since we have verified that the codes T66 and T66' have an automorphism group of order 5, our codes are inequivalent to the two codes.

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- ¹ Department of Mathematical Sciences, Yamagata University, Yamagata 990–8560, Japan

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² Department of Informatics, University of Bergen, Thormøhlensgate 55, N-5008, Bergen, NORWAY