

# Some LP bounds for covering radius of spherical designs

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# Outline of the talk

- Spherical designs (SphD)
- Covering radius of spherical designs
- Gegenbauer polynomials; another characterization of SphD
- LP bounds (obtained by Fazekas and Levenshtein) for covering radius of spherical designs
- Some notations for the structure of spherical designs
- A signed measures and corresponding orthogonal polynomials
- Properties of the polynomials  $P_i^{0,\ell}(t)$
- Improving the Fazekas-Levenshtein (FL) bound for even strengths
- Case:  $\ell \leq t_1(y)$ ; Case:  $t_1(y) \in [-1, \ell]$
- A procedure for finding new lower bounds
- Upper bounds
- Examples for 4-designs
- Future work
- References

## Spherical designs (1)

Spherical designs were introduced in 1977 by Delsarte-Goethals-Seidel [3].

### Definition (1)

A *spherical  $\tau$ -design*  $C \subset \mathbb{S}^{n-1}$  is a finite subset of  $\mathbb{S}^{n-1}$  such that

$$\frac{1}{\mu(\mathbb{S}^{n-1})} \int_{\mathbb{S}^{n-1}} f(x) d\mu(x) = \frac{1}{|C|} \sum_{x \in C} f(x)$$

( $\mu(x)$  is the Lebesgue measure) holds for all polynomials  $f(x) = f(x_1, x_2, \dots, x_n)$  of degree at most  $\tau$  (i.e. the average of  $f$  over the set is equal to the average of  $f$  over  $\mathbb{S}^{n-1}$ ).

The maximal possible  $\tau = \tau(C)$  is called strength of  $C$ .

# Covering radius of spherical designs

## Definition (2)

Let  $C \subset \mathbb{S}^{n-1}$  be a finite set (spherical design in our applications). For a fixed point  $y \in \mathbb{S}^{n-1}$  the *distance between  $y$  and  $C$*  is defined in the usual way by

$$d(y, C) := \min\{d(y, x) : x \in C\}.$$

Then the covering radius of  $C$  is

$$r(C) := \max\{d(y, C) : y \in \mathbb{S}^{n-1}\}.$$

We consider the equivalent quantity

$$\rho(C) := 1 - \frac{r^2(C)}{2} = \min_{y \in \mathbb{S}^{n-1}} \max_{x \in C} \{\langle x, y \rangle\}.$$

# Gegenbauer polynomials

## Definition (3)

By  $P_i^{(n)}(t)$ ,  $i = 0, 1, \dots$ , we denote the *Gegenbauer polynomials* normalized by  $P_i^{(n)}(1) = 1$ , which satisfy the following three-term recurrence relation

$$(i + n - 2) P_{i+1}^{(n)}(t) = (2i + n - 2) t P_i^{(n)}(t) - i P_{i-1}^{(n)}(t) \text{ for } i \geq 1,$$

where  $P_0^{(n)}(t) := 1$  and  $P_1^{(n)}(t) := t$ .

In the standard Jacobi polynomial notation, we have that

$$P_i^{(n)}(t) = \frac{P_i^{((n-3)/2, (n-3)/2)}(t)}{P_i^{((n-3)/2, (n-3)/2)}(1)}.$$

## Spherical designs (2)

### Theorem (4)

[3] A code  $C \subset \mathbb{S}^{n-1}$  is a spherical  $\tau$ -design if and only if for any point  $y \in \mathbb{S}^{n-1}$  and any real polynomial  $f(t)$  of degree at most  $\tau$ , the equality

$$\sum_{x \in C} f(\langle x, y \rangle) = f_0 |C| \quad (1)$$

holds, where  $f_0$  is the first coefficient in the Gegenbauer expansion  $f(t) = \sum_{i=0}^k f_i P_i^{(n)}(t)$ .

## Linear programming (LP) bounds - Fazekas-Levenshtein bounds for covering radius of spherical designs

Linear programming bounds for covering radius of spherical designs were obtained by Fazekas and Levenshtein [2, Theorem 2] in the more general setting of polynomial metric spaces.

They prove that if  $C$  is a  $(2k - 1 + e)$ -design,  $e \in \{0, 1\}$ , then

$$\rho(C) \geq t_{FL} = t_k^{0,e}, \quad (2)$$

where  $t_k^{0,e}$  is the largest zero of the Jacobi polynomial  $P_k^{(\alpha,\beta)}(t)$ ,  $\alpha = \frac{n-3}{2}$ ,  $\beta = \frac{n-3}{2} + e$ .

For example, (2) gives  $\rho(C) \geq (1 + \sqrt{n+3})/(n+2)$  for every spherical 4-design of  $M \geq n(n+3)/2$  points.

Note that the Fazekas-Levenshtein bound does not depend on the cardinality of the designs under consideration.

# Problems

## Problem (1)

*To improve the lower bounds on the covering radius  $\rho(C)$  of spherical designs with fixed dimension  $n$ , even strength  $2k$ , and cardinality  $|C|$ .*

## Problem (2)

*To obtain the upper bounds on the covering radius  $\rho(C)$  of spherical designs with fixed dimension  $n$ , even strength  $2k$ , and cardinality  $|C|$ .*



## Delsarte-Goethals-Seidel bound

For fixed dimension  $n \geq 2$  and strength  $\tau \geq 1$  the minimum cardinality of a spherical  $\tau$ -design  $C \subset \mathbb{S}^{n-1}$  is bounded from below by Delsarte-Goethals-Seidel as follows:

$$|C| \geq D(n, \tau) := \binom{n+k-2+e}{n-1} + \binom{n+k-2}{n-1}, \quad (3)$$

where  $\tau = 2k - 1 + e$ ,  $e \in \{0, 1\}$ . This bound (3) is rarely attained. Improvements (Boyvalenkov, Yudin, Nikova-Nikov); It is still unknown, if there exist spherical 4-designs of 10 points on  $\mathbb{S}^2$ . (Bondarenko, Radchenko and Viazovska) For fixed dimension  $n$  and strength  $\tau$  there exist spherical  $\tau$ -designs on  $\mathbb{S}^{n-1}$  for any cardinality  $N \geq C_n \tau^{n-1}$ , where the constant  $C_n$  depends on the dimension  $n$  only.

## Some notations for the structure of spherical designs

Let  $C \in \mathbb{S}^{n-1}$  be a spherical design. For an arbitrary point  $y \in \mathbb{S}^{n-1}$ , consider the (multi)set

$$I(y) = \{\langle x, y \rangle : x \in C\} = \{t_1(y), t_2(y), \dots, t_{|C|}(y)\},$$

where we assume that  $I(y)$  is ordered by

$$-1 \leq t_1(y) \leq t_2(y) \leq \dots \leq t_{|C|}(y) \leq 1, \quad (t_{|C|}(y) = 1 \Leftrightarrow y \in C).$$

In what follows we always assume that  $y$  is a point on  $\mathbb{S}^{n-1}$  where the covering radius is realized, in particular  $t_{|C|}(y) = \rho(C)$ .

### Lemma (5)

*If  $y$  is a point on  $\mathbb{S}^{n-1}$  where the covering radius is realized, then*

$$t_{|C|}(y) = t_{|C|-1}(y) = \dots = t_{|C|-n+1}(y) = \rho(C).$$

## A signed measure $\Rightarrow$ orthogonal polynomials (1)

The measure of orthogonality of Gegenbauer polynomials  $P_i^{(n)}(t)$  is

$$d\mu(t) := c_n(1 - t^2)^{\frac{n-3}{2}} dt, \quad t \in [-1, 1], \quad c_n := \Gamma\left(\frac{n}{2}\right) / \sqrt{\pi} \Gamma\left(\frac{n-1}{2}\right).$$

We need also measures which are positive definite up to certain degree considered by Cohn and Kumar (2006).

### Definition (6)

A signed Borel measure  $\nu$  on  $\mathbb{R}$  for which all polynomials are integrable is called *positive definite up to degree  $m$*  if  $\int p^2(t) d\nu(t) > 0$  for all real nonzero polynomials  $p(t)$  of degree at most  $m$ .

## A signed measure $\Rightarrow$ orthogonal polynomials (2)

It was proved (B.-Dragnev-Hardin-Saff-S., 2019) that the signed measure

$$d\mu_\ell(t) := c_{n,\ell}(t - \ell)d\mu(t), \quad t \in [-1, 1], \quad c_{n,\ell} := -1/\ell,$$

is positive definite up to degree  $k - 1$  provided that  $\ell < t_{k,1}$ , where  $t_{k,1}$  is the smallest zero of the Gegenbauer polynomial  $P_k^{(n)}(t)$ .

This implies the existence of a finite sequence of polynomials  $\{P_i^{0,\ell}(t)\}_{i=0}^k$  which are orthogonal with respect to  $d\mu_\ell(t)$  and normalized by  $P_i^{0,\ell}(1) = 1$ . The uniqueness of these polynomials allows us to write explicitly

$$P_i^{0,\ell}(t) = \frac{T_i(t, \ell)}{T_i(1, \ell)} = \frac{(1 - \ell) \left( P_{i+1}^{(n)}(t) - P_i^{(n)}(t)P_{i+1}^{(n)}(\ell)/P_i^{(n)}(\ell) \right)}{(t - \ell) \left( 1 - P_{i+1}^{(n)}(\ell)/P_i^{(n)}(\ell) \right)} \quad (4)$$

The boundary case  $\ell = -1$  leads to polynomials which Levenshtein denoted by  $P_i^{0,1}(t)$  - the (normalized) Jacobi polynomials with parameters  $(\alpha, \beta) = ((n - 3)/2, (n - 1)/2)$ .

## Properties of the polynomials $P_i^{0,\ell}(t)$ - interlacing of roots

Let  $\tau = 2k$ . Denote by  $t_{i,1} < t_{i,2} < \dots < t_{i,i}$  the zeros of  $P_i^{(n)}(t)$  and  $t_{i,1}^{0,\ell} < t_{i,2}^{0,\ell} < \dots < t_{i,i}^{0,\ell}$  the zeros of  $P_i^{0,\ell}(t)$ .

### Theorem (7)

Let  $\ell$  and  $k$  be such that  $t_{k+1,1} < \ell < t_{k,1}$  and  $P_{k+1}^{(n)}(\ell)/P_k^{(n)}(\ell) < 1$ . Then

$$P_i^{0,\ell}(t) = \frac{T_i(t, \ell)}{T_i(1, \ell)}, \quad i = 0, 1, \dots, k, \quad (5)$$

where the leading coefficient of  $P_i^{0,\ell}$  is  $m_i^{0,\ell} > 0$ . All zeros  $\{t_{i,j}^{0,\ell}\}_{j=1}^i$  of  $P_i^{0,\ell}(t)$  are in the interval  $[\ell, 1]$  and the interlacing rules

$$\begin{aligned} t_{i,j}^{0,\ell} &\in (t_{i,j}, t_{i+1,j+1}), \quad i = 1, \dots, k-1, j = 1, \dots, i; \\ t_{k,j}^{0,\ell} &\in (t_{k+1,j+1}, t_{k,j+1}), \quad j = 1, \dots, k-1, \quad t_{k,k}^{0,\ell} \in (t_{k+1,k+1}, 1), \end{aligned} \quad (6)$$

hold.

## Properties of the polynomials $P_i^{0,\ell}(t)$ - a quadrature formula

We denote by  $L_i(t)$ ,  $i = 0, 1, \dots, k$ , the Lagrange basic polynomials generated by the nodes  $\ell < t_{k,1}^{0,\ell} < t_{k,2}^{0,\ell} < \dots < t_{k,k}^{0,\ell}$  and set

$$\theta_i := \int_{-1}^1 L_i(t) d\mu(t), \quad i = 0, 1, \dots, k.$$

### Theorem (8)

Let  $t_{k,1}^{0,\ell} < t_{k,2}^{0,\ell} < \dots < t_{k,k}^{0,\ell}$  be the zeros of the polynomial  $P_k^{0,\ell}(t)$ . Then the quadrature formula

$$f_0 = \int_{-1}^1 f(t) d\mu(t) = \theta_0 f(\ell) + \sum_{i=1}^k \theta_i f(t_{k,i}^{0,\ell}) \quad (7)$$

is exact for all polynomials of degree at most  $2k$  and has positive weights  $\theta_i > 0$ ,  $i = 0, 1, \dots, k$ .

## Improving the FL bound for even strengths

$$\text{Case: } -1 < \ell \leq t_1(y) \quad (1)$$

Let  $C \subset \mathbb{S}^{n-1}$  be a spherical  $2k$ -design of cardinality  $|C| > D(n, 2k)$ . Let  $y \in \mathbb{S}^{n-1}$  be a point which realizes the covering radius of  $C$ .

Lower bounds  $-1 < \ell \leq t_1(y)$  imply improvements of the  $t_{FL}$ .

### Theorem (9)

*Let  $C \subset \mathbb{S}^{n-1}$  be a spherical  $2k$ -design and  $y \in \mathbb{S}^{n-1}$  be a point which realizes the covering radius of  $C$ . If  $\ell \leq t_1(y)$ , then*

$$\rho(C) \geq t_{k,k}^{0,\ell}.$$

In the boundary case  $\ell = -1$  this theorem gives the FASEKAS-LEVENSHTEIN bound  $\rho(C) \geq t_{FL} = t_k^{0,1}$ . Therefore, we have improvement of the FASEKAS-LEVENSHTEIN bound whenever it is known (or it is presumed) that  $\ell \leq t_1(y)$  for a point  $y$  where the covering radius is realized.

Case:  $-1 < \ell \leq t_1(y)$       Some lower bounds  $t_{k,k}^{0,\ell}$       (2)

Dimension $n$	Cardinality $ C $	Strength $\tau = 2k$	$\ell$	FL-lower bound $\rho(C) \geq t_k^{0,1}$	New lower bound if $\ell \leq t_1(y)$ $\rho(C) \geq t_{k,k}^{0,\ell}$
3	10	4	-0.97	0.689897	0.694892
3	10	4	-0.95	0.689897	0.698664
3	10	4	-0.9	0.689897	0.710257
4	15	4	-0.97	0.607625	0.611772
4	15	4	-0.95	0.607625	0.614815
4	15	4	-0.9	0.607625	0.623682
3	17	6	-0.97	0.822824	0.825859
3	17	6	-0.95	0.822824	0.828450
3	17	6	-0.9	0.822824	0.839165



## Improving the FL bound for even strengths

$$\text{Case: } t_1(y) \in [-1, \ell] \quad (1)$$

This case is more subtle.

We will have to optimize in two classes of real polynomials. We consider

$$A(n, k, \ell) := \{f(t) = A^2(t) : \deg(f) = 2k, A(t) \text{ has } k \text{ real zeros in } [\ell, t_{FL}]\}.$$

Similarly, we use polynomials from the set

$$B(n, k, s) := \{g(t) = (t + 1)B^2(t)(t - s) : \deg(g) = 2k,$$

$$B(t) \text{ has } k - 1 \text{ real zeros in } [-1, s]\},$$

where the parameter  $s$  (close to  $t_{FL}$ ) will be chosen in advance.

Case:  $t_1(y) \in [-1, \ell]$  (2)

The next lemma sets an auxiliary parameter  $m(C)$  after optimization in the set  $A(n, \tau, \ell)$ .

Lemma (10)

Let  $f(t) \in A(n, k, \ell)$  and the positive integer  $m$  be such that

$$f_0|C| < f(\ell) + (m+1)f(t_{FL}). \quad (8)$$

Then  $t_{|C|-m}(y) < t_{FL}$ .

We define

$$m(C) := \min\{m : \exists f \in A(n, k, \ell) \text{ such that } f_0|C| < f(\ell) + (m+1)f(t_{FL})\}.$$

This lemma implies that

$$m(C) \geq n.$$

Case:  $t_1(y) \in [-1, \ell]$  (3)

### Example (11)

We have  $m(C) = n = 3$  for  $(n, \tau, |C|) = (3, 4, 10)$ , the first case, where the existence/nonexistence of spherical 4-designs is undecided. Similarly, for  $(n, \tau, |C|) = (4, 4, 15)$  we have  $m(C) = n + 1 = 5$ .

### Lemma (12)

Let  $f(t) \in A(n, k, \ell)$  be such that  $f_0|C| < f(\ell) + (m(C) + 1)f(t_{FL})$ . Then  $t_{|C|-m(C)}(y) \leq s$ , where  $s$  is the largest root of the equation

$$f_0|C| - f(\ell) = (m(C) + 1)f(t).$$

$$\text{Case: } t_1(y) \in [-1, \ell] \quad (4)$$

The previous two Lemmas 10 and 12 imply that  $t_{|C|-m(C)}(y) \leq s < t_{FL}$ . We utilize this in a second optimization dealing with the location of  $t_{FL}$  between two inner products from  $I(y)$ .

We have

$$t_{|C|-m(C)}(y) \leq s < t_{FL} \leq t_{|C|-n+1}(y) = \rho(C).$$

Therefore, there exist  $j \in \{0, 1, \dots, m(C) - n\}$  such that

$$t_{|C|-m(C)+j}(y) < t_{FL} \leq t_{|C|-m(C)+j+1}(y). \quad (9)$$

This clarification of the location of  $t_{FL}$  with respect to the points of  $I(y)$  allows more precise estimations.

Case:  $t_1(y) \in [-1, \ell]$  (5)

Lemma (13)

If  $g(t) \in B(n, k, s)$ , then  $\rho(C) \geq m_{\ell, s}^{(j)}$ , where  $m_{\ell, s}^{(j)}$  is the largest root of the equation

$$jg(t_{FL}) + (m(C) - j)g(t) = g_0|C|. \quad (10)$$

Lemma (14)

Let  $f(t) \in A(n, k, \ell)$  be such that  $f_0|C| < f(\ell) + (m(C) + 1)f(t_{FL})$ . Then  $t_{|C|-m(C)}(y) \leq s^{(j)}$ , where  $s^{(j)}$  is the largest root of the equation

$$f_0|C| = (j + 1)f(t) + f(\ell) + (m(C) - j)f(t_{FL}). \quad (11)$$

## A procedure for finding new lower bounds (1)

In the case  $t_1(y) \in [-1, \ell]$ , for each fixed  $j \in \{0, 1, \dots, m(C) - n\}$ , we can start an iterative procedure with previous Lemmas 13 and 14 for obtaining consecutive improvements of  $s^{(j)}$  and  $m_{\ell, s}^{(j)}$ .

This procedure may converge to some bounds or may be divergent which will mean nonexistence of designs with the corresponding parameters (dimension, strength, and cardinality).

The better bound  $\rho(C) \geq t_k^{0, \ell}$  when  $t_1(y) \geq \ell$  allows starting a similar procedure with analogs of the above Lemmas 13 and 14 as the only difference will be the absence of  $\ell$ .

## A procedure for finding new lower bounds (2)

### Example (15)

Considering again  $(n, \tau, |C|) = (3, 4, 10)$

(recall that  $m(C) = n = 3$  in this case, i.e.  $j = 0$  only),

we obtain for  $\ell = -0.97$  that  $\rho(C) \geq 0.724753$  if  $t_1(y) \in [-1, -0.97]$

and  $\rho(C) \geq 0.728787$  if  $t_1(y) \geq -0.97$ .

Therefore, we have  $\rho(C) \geq 0.724753$  in the worst case.

Similarly, for  $(n, \tau, |C|) = (4, 4, 15)$

(note that now  $m(C) = 5$ , i.e.  $j = 0, 1$ ),

we obtain for  $\ell = -0.97$  that  $\rho(C) \geq 0.625572$  if  $t_1(y) \in [-1, -0.97]$

and  $\rho(C) \geq 0.627354$  if  $t_1(y) \geq -0.97$  for  $j = 0$ ;

$\rho(C) \geq 0.616854$  if  $t_1(y) \in [-1, -0.97]$

and  $\rho(C) \geq 0.619259$  if  $t_1(y) \geq -0.97$  for  $j = 1$ .

Summarizing, we conclude that  $\rho(C) \geq 0.616854$  in the worst case.

## General upper bounds (1)

We write (1) for  $y$ ,  $C$  and  $f(t)$ ,  $\deg(f) \leq \tau(C)$ , as

$$nf(\rho(C)) + \sum_{i=1}^{|C|-n} f(t_i(y)) = f_0|C|. \quad (12)$$

The identity (12) provides upper bounds for  $\rho(C)$  as follows.

### Theorem (16)

(Linear programming upper bounds of the covering radius of spherical designs) *Let  $f(t)$ ,  $\deg(f) \leq \tau$ , be a real polynomial which is nonnegative in  $[-1, t_{FL}]$  and increasing in  $[t_{FL}, 1]$ . Then for every  $\tau$ -design  $C \subset \mathbb{S}^{n-1}$  we have*

$$\rho(C) \leq m_u,$$

where  $m_u$  is the largest root of the equation  $nf(t) = f_0|C|$ .



## General upper bounds (2)

The following theorem shows which kind of extremal polynomials should be investigated.

### Theorem (17)

*The best polynomials for use in Theorem 16 are  $f(t) = (t + 1)^e A^2(t)$ , where  $\tau = 2k - e$ ,  $e \in \{0, 1\}$ ,  $\deg(A) = k - e$  and  $A(t)$  has  $k - e$  zeros in  $[-1, t_{FL}]$ .*

## Upper bounds for spherical 4-designs

We now find the optimal polynomials in the above Theorem for  $\tau = 4$ .

### Theorem (18)

If  $C \subset \mathbb{S}^{n-1}$  is a spherical 4-design, then

$$\rho(C) \leq u(a_0, b_0),$$

where the function  $u(a, b)$  and the (optimal) parameters  $a_0$  and  $b_0$  are defined in the proof.

### Example (19)

Looking again in the first open case  $(n, \tau, |C|) = (3, 4, 10)$ , we obtain (for

$b_0 = \frac{\sqrt{69}-7}{30}$  and  $a_0 = \frac{(3+\sqrt{69})\sqrt{45+10\sqrt{69}}}{150}$ ) the upper bound  
 $\rho(C) \leq u(a_0, b_0) \approx 0.754443$ .

## Examples for 4-designs

Dimension $n$	Cardinality $ C $	$m(C)$	FL-lower bound $\rho(C) \geq t_2^{0,1}$	New lower bound $\ell = -0.97$	New upper bound
3	10	3	0.689897	0.724753	0.7545
3	11	4	0.689897	0.694717	0.7794
4	15	5	0.607625	0.616854	0.6918
4	16	5	0.607625	0.610537	0.7072
5	21	7	0.546918	0.550012	0.6503
5	22	8	0.546918	0.548132	0.6604
6	28	10	0.500000	0.501717	0.6198
6	29	10	0.500000	0.501288	0.6269
7	36	13	0.462475	0.463455	0.5960
7	37	13	0.462475	0.462961	0.6012

## Upper bounds for antipodal 3- and 5-designs (1)

A spherical design  $C$  is called antipodal if  $C = -C$ .

The set  $I(y)$  is symmetric for antipodal designs;

We have

$$t_i = t_{|C|-i+1}, \quad i = 1, 2, \dots, n.$$

In Lemma 5 and the equation in Theorem (16) becomes

$$2nf(t) = f_0|C|.$$

The upper bounds for antipodal designs are easier for  $\tau = 3$  and  $\tau = 5$ .

We are able to obtain an explicit bound in these cases for all dimensions and cardinalities.

## Upper bounds for antipodal 3- and 5-designs (2)

### Theorem (20)

If  $C$  is an antipodal 3-design, then

$$t_{FL} = \frac{1}{\sqrt{n}} \leq t_c \leq \frac{1}{n} \sqrt{\frac{|C|}{2}}.$$

### Theorem (21)




If  $C$  is an antipodal 5-design, then

$$t_{FL} = \left(\frac{3}{n+2}\right)^{1/2} \leq t_c \leq \left(\frac{1}{n} + \frac{1}{n} \sqrt{\frac{(n-1)(|C|-2n)}{n(n+2)}}\right)^{1/2}.$$

## Future work

- To obtain an explicit bounds for covering radius of spherical designs with  $\tau = 6, 8$ , etc.
- To obtain bounds for covering radius of spherical designs with odd strengths.

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