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DIRAC TYPE CONDITION AND HAMILTONIAN GRAPHS

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ABSTRACT. In 1952, Dirac introduced the degree type condition and proved that if G is a connected graph of order $n \ge 3$ such that its minimum degree satisfies $\delta(G) \ge n/2$, then G is Hamiltonian. In this paper we investigate a further condition and prove that if G is a connected graph of order $n \ge 3$ such that $\delta(G) \ge (n-2)/2$, then G is Hamiltonian or G belongs to four classes of well-structured exceptional graphs.

1. Introduction. We consider only finite undirected graphs without loops or multiple edges. For a graph G, let V(G) be the vertex set and let E(G)be the edge set of G. The complete graph of order n is denoted by K_n and the empty graph of order n is denoted by \overline{K}_n . For two vertices u and v, let d(u, v)be the length of the shortest path between vertices u and v in G, i.e., d(u, v) is the distance between u and v. The minimum degree of the graph G is denoted by $\delta(G)$. For a subgraph H of the graph G and a subset S of V(G), let $N_H(S)$ be the set of vertices in H that are adjacent to some vertex in S and let the

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cardinality of $N_H(S)$ be $d_H(S)$. Furthermore, we denote by G - H the subgraph of G induced by V(G) - V(H). For each integer $m \ge 3$, let $C_m = x_1 x_2 \cdots x_m x_1$ denote a cycle of order m and let P be a path of a section of C_m (if |V(P)| = m, then $P = C_m$). Define

$$N_P^+(u) = \{x_{i+1} : x_i \in N_P(u)\}, N_P^-(u) = \{x_{i-1} : x_i \in N_P(u)\}, N_P^-(u)$$

and define $N_P^{\pm}(u) = N_P^{+}(u) \cup N_P^{-}(u)$, where subscripts are taken modulo m.

If no ambiguity can arise we sometimes write N(u) instead of $N_G(u)$, δ instead of $\delta(G)$, etc. Other notations can be found in [1, 3].

In 1952, Dirac proved the following well-known result on Hamiltonian graphs.

Theorem 1 (Dirac [2]). If G is a connected graph of order $n \ge 3$ and $\delta(G) \ge n/2$, then G is Hamiltonian.

Recently we also have investigated some Hamiltonian graphs under other sufficient conditions such as neighborhood union conditions [4, 5].

In this paper, our purpose is to present the following result, which extends the above Theorem 1.

Theorem 2. If G is a connected graph of order $n \ge 3$ and $\delta(G) \ge (n-2)/2$, then G is Hamiltonian or G has one of the following four types:

$$G_{(n-2)/2} \vee (\overline{K}_{(n-2)/2} \cup K_2), \quad G_{(n-2)/2} \vee \overline{K}_{(n+2)/2},$$

 $G_{(n-1)/2} \vee \overline{K}_{(n+1)/2}, \quad w : (K_h \cup K_t^-).$

Here h is an integer, G_h denotes an arbitrary graph of order h, and \overline{K}_h is the empty graph with h vertices and without edges. The join operator $A \vee B$ of two graphs A and B is the graph constructed from A and B by adding all edges joining the vertices of A and the vertices of B. The graph $A \cup B$ denotes the disjoint union of the graphs A and B. The graph $w : (K_h \cup K_t^-)$ is defined by the properties: K_t^- is a graph of order t and minimum degree $\delta(K_t^-) \geq t - 2$, w is a cut vertex which is adjacent to at least (n-2)/2 vertices of the disjoint graphs K_h and K_t^- , $(n-2)/2 \leq h \leq t \leq n/2$.

Corollary 3. If G is a 2-connected graph of order $n \ge 3$ and $\delta(G) \ge (n-2)/2$, then G is Hamiltonian or

$$G \in \{G_{(n-2)/2} \lor (\overline{K}_{(n-2)/2} \cup K_2), G_{(n-2)/2} \lor \overline{K}_{(n+2)/2}, G_{(n-1)/2} \lor \overline{K}_{(n+1)/2}.$$

2. The proof of main Theorem.

Proof of Theorem 2. Assume that G satisfies the condition of Theorem 2 and it is not Hamiltonian. Then let $C_m = x_1 x_2 \cdots x_m x_1$ be the longest cycle of G and let H be a component of $G - C_m$. We consider two cases.

Case 1. G is 2-connected.

In this case, there must exist vertices $u, v \in V(H)$ such that $x_i \in N_{C_m}(u)$ and $x_j \in N_{C_m}(v)$. (If |V(H)| = 1, then u = v.) Now we claim that $d(x_{i+1}) + d(x_{j+1}) \leq n - |V(H)|$. Otherwise, if the claim is false, let P and R denote, respectively, the path $x_{i+1}x_{i+2}\cdots x_j$ of C_m and the path $x_{j+1}x_{j+2}\cdots x_i$ of C_m . Then, clearly, none of $N_P^+(x_{j+1}) \cup N_R^-(x_{j+1})$ are adjacent to x_{i+1} . (For example, if $x_k \in N_P^+(x_{j+1})$ is adjacent to x_{i+1} , let T be a path in H which has two endvertices adjacent to x_i and x_j , respectively. Then the cycle

$$x_i T x_j x_{j-1} \cdots x_k x_{k-1} \cdots x_{i+1} x_{k-1} x_{j+1} x_{j+2} \cdots x_i$$

is longer than C_m , a contradiction.) Clearly $|N_P^+(x_{j+1}) \cup N_R^-(x_{j+1})| = |N_{C_m}(x_{j+1})| - |\{x_{j+1}\}|$. Also, none of $N_{G-C_m}(x_{j+1})$ are adjacent to x_{i+1} , and both x_{i+1}, x_{j+1} are not adjacent to any vertex of $\{x_{i+1}\} \cup V(H)$. Hence we can check that

$$\begin{aligned} d(x_{i+1}) &\leq |V(G)| - |N_P^+(x_{j+1}) \cup N_R^-(x_{j+1})| - |N_{G-C_m}(x_{j+1})| - |\{x_{i+1}\} \cup V(H)| \\ &\leq |V(G)| - (|N_{C_m}(x_{j+1})| - |\{x_{j+1}\}|) - |N_{G-C_m}(x_{j+1})| - |\{x_{i+1}\} \cup V(H)| \\ &\leq n - (|N(x_{j+1})| - |\{x_{j+1}\}|) - |\{x_{i+1}\}| - |V(H)|, \end{aligned}$$

and this implies that

(1)
$$d(x_{i+1}) + d(x_{j+1}) \le n - |V(H)|.$$

On the other hand, by the condition of Theorem 2 we have $d(x_{i+1}) + d(x_{j+1}) \ge n-2$. Together with the inequality (1), we have $|V(H)| \le 2$. Now we consider the subcases |V(H)| = 1 and |V(H)| = 2.

Subcase 1.1. |V(H)| = 2.

In this case, if $u \in V(H)$, then we have $|j - i| \geq 3$ for all pairs $x_i, x_j \in N_{C_m}(H)$ such that $\{x_{i+1}, x_{i+2}, \cdots, x_{j-1}\} \cap N(u) = \emptyset$. (Otherwise, if $|j - i| \leq 2$ for some $x_i, x_j \in N_{C_m}(H)$, by |V(H)| = 2, it is easy to construct a cycle longer than C_m , a contradiction.) Thus, we can check that

$$|N(u)| \le |V(C_m)|/3 + |V(H) \setminus \{u\}| \le (n-2)/3 + 1,$$

i.e., $d(u) \leq (n-2)/3+1$. Clearly, $n \geq |V(C_m)| + |V(H)| \geq 8$, so $d(u) \leq (n-3)/2$, and this contradicts the assumption of Theorem 2 that $d(u) \geq (n-2)/2$.

Subcase 1.2. |V(H)| = 1.

In this case, if $u \in V(H)$, then we have $|j - i| \ge 2$ for all pairs $x_i, x_j \in N_{C_m}(H)$ with $\{x_{i+1}, x_{i+2}, \dots, x_{j-1}\} \cap N(u) = \emptyset$, so $|N(u)| \le |V(C_m)|/2 \le (n-1)/2$. Together with the assumption of Theorem 2 that $d(u) \ge (n-2)/2$, we have $(n-2)/2 \le d(u) \le (n-1)/2$. Then we consider the following subcases.

Subcase 1.2.1. d(u) = (n-2)/2.

Then, since C_m is the longest cycle of G, the vertex u is not adjacent to two consecutive x_i, x_{i+1} on C_m . By $d(u) = (n-2)/2, |V(C_m)| \ge 2d(u) \ge n-2$, so $|V(G - C_m)| \le 2$.

First, let $|V(G - C_m)| = 1$. In this case, since G does not have a Hamiltonian cycle C_n and d(u) = (n-2)/2, it is easy to obtain $N(u) = \{x_i, x_{i+3}, x_{i+5}, x_{i+7}, \dots, x_{i-2}\}$ on C_m , i.e., there exists only one pair of consecutive neighbor vertices x_i, x_{i+3} of u on C_n such that |(i+3)-i| = 3 and for all other two consecutive neighbor vertices x_{i+k}, x_{i+h} of u with $\{x_{i+k+1}, x_{i+k+2}, \dots, x_{i+h-1}\} \cap N(u) = \emptyset$, we have |(i+k) - (i+h)| = 2.

In this case, since G does not have a Hamiltonian cycle C_n , we derive that

$$G \in G_{(n-2)/2} \lor (\overline{K}_{(n-2)/2} \cup K_2),$$

where

$$V(G_{(n-2)/2}) = \{x_i, x_{i+3}, x_{i+5}, x_{i+7}, \cdots x_{i-2}\},\$$

$$\overline{K}_{(n-2)/2} = \{x_{i+4}, x_{i+6}, \cdots x_{i-1}, u\}, \quad K_2 = \{x_{i+1}, x_{i+2}\}.$$

Now, let $|V(G - C_m)| = 2$ and let $v \in V(G - C_m - u)$. Since G does not have a Hamiltonian cycle C_n and d(u) = (n-2)/2, it is easy to obtain $N(u) = N(v) = \{x_i, x_{i+1}, \dots, x_{i-2}\}$. This implies

$$G \in G_{(n-2)/2} \vee \overline{K}_{(n+2)/2},$$

where

$$V(G_{(n-2)/2}) = \{x_i, x_{i+2}, \cdots, x_{i-2}\}, \overline{K}_{(n+2)/2} = \{x_{i+1}, x_{i+3}, \cdots, x_{i-1}, u, v\}.$$

Subcase 1.2.2. d(u) = (n-1)/2.

In this case, since C_m is the longest cycle of G, the vertex u is not adjacent to two consecutive x_i, x_{i+1} on C_m , so $|V(C_m)| \ge 2d(u) \ge n-1$. This implies $|V(G - C_m)| = 1$. In this case, it is easy to obtain $N(u) = \{x_i, x_{i+2}, \dots, x_{i-2}\}$ on C_m , so we have

$$G \in G_{(n-1)/2} \lor \overline{K}_{(n+1)/2},$$

where

$$V(G_{(n-1)/2}) = \{x_i, x_{i+2}, \cdots, x_{i-2}\}, \overline{K}_{(n+1)/2} = \{x_{i+1}, x_{i+3}, \cdots, x_{i-1}, u\}.$$

Case 2. The connectivity of G is 1.

In this case, let w be a cut vertex of G. Then, since $d(x) \ge (n-2)/2$ for each vertex x in G, we have that G - w has two components. (Otherwise, let G - w have at least three components H_1 , H_2 and H_3 , and let $|V(H_1)| =$ $\min\{|V(H_1)|, |V(H_2)|, |V(H_3)|\}$. Then we have $d(y) \le n/3 - |\{y\}| + |\{w\}| < (n-2)/2$ for each vertex y in H_1 , a contradiction.) Let H_1, H_2 be the components of G - w, i.e., $G - w = H_1 \cup H_2$. We denote G by $w : (H_1 \cup H_2)$. Clearly, we have

 $(n-2)/2 \le \min\{|V(H_1)|, |V(H_2)|\} \le \max\{|V(H_1)|, |V(H_2)|\} \le n/2.$

Without loss of generality we may assume that

$$|V(H_1)| = h = \min\{|V(H_1)|, |V(H_2)|\}$$
$$|V(H_2)| = t = \max\{|V(H_1)|, |V(H_2)|\}.$$

(i) When h = (n-2)/2, then H_1 is the complete subgraph K_h , the vertex w is adjacent to all vertices of H_1 , and each vertex of H_2 is not adjacent to at least one vertex of H_2 . (ii) When h = (n-1)/2, then H_1 and H_2 are both complete subgraphs. However, we can write $w : (H_1 \cup H_2) = w : (K_h \cup K_t^-)$, where $(n-2)/2 \le h \le t \le n/2$, and each vertex of K_t^- is not adjacent to at least one vertex of K_t^- .

Therefore, the proof is complete. \Box

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