

INNER ESTIMATION  
OF LINEAR PARAMETRIC  $AE$ -SOLUTION SETS

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**Abstract**

We consider linear algebraic systems  $A(p)x = b(p)$ , where the elements of the matrix and the right-hand side vector are linear functions of uncertain parameters varying within given intervals. For the parametric  $AE$ -solution set defined for two disjoint index sets  $\mathcal{A}, E$  by  $\Sigma_{AE}^p := \{x \in \mathbb{R}^n \mid (\forall p_{\mathcal{A}} \in [p_{\mathcal{A}}]) (\exists p_{\mathcal{E}} \in [p_{\mathcal{E}}]) (A(p)x = b(p))\}$ , an inner interval estimation  $[y] \subseteq \Sigma_{AE}^p$  is sought. In the special case when  $\Sigma_{AE}^p$  has linear shape, this paper gives parametric generalization of the so-called centred approach, which was developed so far only for tolerable solution sets.

**Key words:** interval linear equations, dependent data,  $AE$ -solution sets, inner inclusion

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**1. Introduction.** Consider linear algebraic systems having linear uncertainty structure

$$(1) \quad \begin{aligned} A(p)x &= b(p), \\ A(p) &:= A_0 + \sum_{k=1}^K p_k A_k, \quad b(p) := b_0 + \sum_{k=1}^K p_k b_k, \end{aligned}$$

where  $A_k \in \mathbb{R}^{m \times n}$ ,  $b_k \in \mathbb{R}^m$ ,  $k = 0, \dots, K$  and the parameters  $p = (p_1, \dots, p_K)^\top$  are considered to be uncertain and varying within given intervals  $\mathbf{p} = (\mathbf{p}_1, \dots, \mathbf{p}_K)^\top$ . Above,  $\mathbb{R}^n$  and  $\mathbb{R}^{m \times n}$  denote the set of real vectors with  $n$  components and the set of real  $m \times n$  matrices, respectively. A real compact interval is  $\mathbf{a} = [\underline{a}, \bar{a}] := \{a \in \mathbb{R} \mid \underline{a} \leq a \leq \bar{a}\}$ . By  $\mathbb{I}\mathbb{R}^n$ ,  $\mathbb{I}\mathbb{R}^{m \times n}$  we denote the sets of interval  $n$ -vectors and

interval  $m \times n$  matrices, respectively. For  $\mathbf{a} = [\underline{a}, \bar{a}]$ , define midpoint  $\check{\mathbf{a}} := (\underline{a} + \bar{a})/2$ , radius  $\hat{\mathbf{a}} := (\bar{a} - \underline{a})/2$  and absolute value (magnitude)  $|\mathbf{a}| := \max\{|\underline{a}|, |\bar{a}|\}$ . These functions are applied to interval vectors and matrices componentwise. Without loss of generality and in order to have a unique representation (1), we assume that  $\hat{\mathbf{p}}_k > 0$  for all  $1 \leq k \leq K$ .

We consider the parametric  $AE$ -solution sets of system (1), which are defined by

$$(2) \quad \Sigma_{AE}^p = \Sigma_{AE}(A(p), b(p), \mathbf{p}) \\ := \{x \in \mathbb{R}^n \mid (\forall p_{\mathcal{A}} \in \mathbf{p}_{\mathcal{A}})(\exists p_{\mathcal{E}} \in \mathbf{p}_{\mathcal{E}})(A(p)x = b(p))\},$$

where  $\mathcal{A}$  and  $\mathcal{E}$  are sets of indexes such that  $\mathcal{A} \cup \mathcal{E} = \{1, \dots, K\}$ ,  $\mathcal{A} \cap \mathcal{E} = \emptyset$ . For a given index set  $\Pi = \{\pi_1, \dots, \pi_k\}$ ,  $p_{\Pi}$  denotes  $(p_{\pi_1}, \dots, p_{\pi_k})$ . Among the  $AE$ -solution sets most studied and of particular practical interest are: the (parametric) united solution set

$$\Sigma_{\text{uni}}(A(p), b(p), \mathbf{p}) := \{x \in \mathbb{R}^n \mid (\exists p \in \mathbf{p})(A(p)x = b(p))\},$$

the (parametric) tolerable solution set

$$\Sigma(A(p_{\mathcal{A}}), b(p_{\mathcal{E}}), \mathbf{p}) := \{x \in \mathbb{R}^n \mid (\forall p_{\mathcal{A}} \in \mathbf{p}_{\mathcal{A}})(\exists p_{\mathcal{E}} \in \mathbf{p}_{\mathcal{E}})(A(p_{\mathcal{A}})x = b(p_{\mathcal{E}}))\}$$

and the (parametric) controllable solution set

$$\Sigma(A(p_{\mathcal{E}}), b(p_{\mathcal{A}}), \mathbf{p}) := \{x \in \mathbb{R}^n \mid (\forall p_{\mathcal{A}} \in \mathbf{p}_{\mathcal{A}})(\exists p_{\mathcal{E}} \in \mathbf{p}_{\mathcal{E}})(A(p_{\mathcal{E}})x = b(p_{\mathcal{A}}))\}.$$

Although the nonparametric  $AE$ -solution sets are studied to some extent (see, e.g., [2, 13] and the references given therein), there are only a few results on the more general case of linear parameter dependency. The so-called ‘‘centred’’ approach for inner estimation of a solution set is developed by NEUMAIER in [5] for the nonparametric united solution set and for the nonparametric tolerable solution set. In [9] this approach is generalized for the parametric tolerable solution set. In this paper we expand the parametric centred approach for parametric  $AE$ -solution sets (in particular the most used: united and controllable solution sets) which have linear shape (boundary). Parametric solution sets with linear shape appear often in the applied domains, see, e.g., [6], ([10], Examples 5.2, 5.3).

## 2. Theoretical background.

**Definition 2.1** ([12]). A parameter  $p_k$ ,  $1 \leq k \leq K$ , is of 1st class if it occurs in only one equation of the system (1).

**Definition 2.2** ([12]). A parameter  $p_k$ ,  $1 \leq k \leq K$ , is of 2nd class if it is involved in more than one equation of the system (1).

The parametric centred approach is based on the explicit description of the parametric  $AE$ -solution set. How to obtain it is discussed in [8]. The description

is explicit if the system involves only 1st class existentially quantified parameters, called  $\mathcal{E}$ -parameters for short, and in some other special cases. In the general case, the description can be obtained by an algorithmic procedure. One has first to eliminate all 2nd class  $\mathcal{E}$ -parameters by the Fourier–Motzkin like parameter elimination procedure. The elimination of every 2nd class  $\mathcal{E}$ -parameter introduces, in general, new characterizing inequalities/inclusions, besides those given in ([12], Theorems 4.1, 4.2). The degree of polynomials involved in the characterizing inequalities/inclusions may be arbitrary high and depends on the 2nd class  $\mathcal{E}$ -parameters involved in the parametric matrix.

**Definition 2.3.** A parametric  $AE$ -solution set is called *linear*, in other words its shape is linear, if the boundary of the solution set consists of parts of hyperplanes.

The inequalities/inclusions describing a *linear* parametric  $AE$ -solution set involve only linear functions on the coordinate variables.

**Theorem 2.1.** *Let  $\Sigma_{AE}^p$  have linear shape. A point  $x \in \mathbb{R}^n$  belongs to  $\Sigma_{AE}^p$ , if and only if*

$$(3) \quad |U(\check{\mathbf{p}})x - v(\check{\mathbf{p}})| \leq \sum_{k=1}^K \delta_k \hat{\mathbf{p}}_k |U_k x - v_k|$$

*equivalently*

$$(4) \quad \sum_{k \in \mathcal{A}} \mathbf{p}_k (U_k x - v_k) \subseteq v_0 - U_0 x + \sum_{k \in \mathcal{E}} \mathbf{p}_k (v_k - U_k x),$$

where  $U_k \in \mathbb{R}^{q \times n}$ ,  $v_k \in \mathbb{R}^q$ ,  $k = 0, \dots, K$ ,  $U(p) = U_0 + \sum_{k=1}^K p_k U_k$ ,  $v(p) = v_0 + \sum_{k=1}^K p_k v_k$ ,  $q \geq m$ , correspond to the explicit description of  $\Sigma_{AE}^p$  and  $\delta_k := \{1 \text{ if } k \in \mathcal{E}, -1 \text{ if } k \in \mathcal{A}\}$ .

**Proof.** Inequalities (3) represent the explicit description of a  $\Sigma_{AE}^p$  with linear shape and follow from the Fourier–Motzkin like parameter elimination procedure. The equivalence between (3) and (4) is proven similarly to ([12], Theorem 3.3).  $\square$

We have  $U_i(p) = A_i(p)$  and  $v_i(p) = b_i(p)$  for  $i = 1, \dots, m$ . In the special case when the definition of  $\Sigma_{AE}^p$  involves only 1st class  $\mathcal{E}$ -parameters,  $U(p) = A(p)$ ,  $v(p) = b(p)$ . The same is true for  $\Sigma_{\text{tol}}(A(p), \mathbf{b}, \mathbf{p})$ . Any  $\Sigma(A(p_{\mathcal{A}}), b(p_{\mathcal{E}}), \mathbf{p})$  also has linear shape independent of the number of 2nd class  $\mathcal{E}$ -parameters involved in the right-hand side vector. Some general sufficient conditions for a parametric  $AE$ -solution set to have linear shape will be given in a separate paper.

Since the methodology we consider requires knowledge of a point which belongs to the interior of the solution set, in what follows we assume that  $\Sigma_{AE}^p$  is not empty and non-degenerate (having a nonempty interior). Denote by  $\text{int } \mathcal{X}$  the topological interior of the set  $\mathcal{X}$  in  $\mathbb{R}^n$  with the standard topology.

We need also the so-called inner interval subtraction, defined for  $\mathbf{a}, \mathbf{b} \in \mathbb{IR}$  by

$$\mathbf{a} \ominus \mathbf{b} := [\underline{a} - \underline{b}, \bar{a} - \bar{b}],$$

where  $\mathbf{a} \ominus \mathbf{b} = \emptyset$  if  $\underline{a} - \underline{b} > \bar{a} - \bar{b}$ . Inner interval subtraction can be represented by the arithmetic operations and functions in the generalized interval space of proper and improper intervals known as KAUCHER interval arithmetic [4] or modal interval arithmetic [1]. In order to facilitate the derivations in the next section we will use some properties of the arithmetic on proper and improper intervals.

The set of *proper* intervals  $\mathbb{IR}$  is extended in [4] by the set  $\{[\underline{a}, \bar{a}] \mid \underline{a}, \bar{a} \in \mathbb{R}, \underline{a} \geq \bar{a}\}$  of *improper* intervals obtaining thus the set  $\mathbb{I}^*\mathbb{R} = \{[\underline{a}, \bar{a}] \mid \underline{a}, \bar{a} \in \mathbb{R}\}$  of generalized intervals. The conventional interval arithmetic and lattice operations, order relations and other functions are isomorphically extended onto the whole set  $\mathbb{I}^*\mathbb{R}$ , [4]. Modal interval analysis [1] imposes a logical-semantic background on generalized intervals (considered there as modal intervals) and allows giving a logical meaning to the interval results. The conventional interval arithmetic can be obtained as a projection of the generalized interval arithmetic on  $\mathbb{IR}$ . An element-to-element symmetry between proper and improper intervals is expressed by the “dual” operator  $\mathbf{dual}(\mathbf{a}) := [\bar{a}, \underline{a}]$  for  $\mathbf{a} = [\underline{a}, \bar{a}] \in \mathbb{I}^*\mathbb{R}$ .  $\mathbf{dual}$  is applied componentwise to vectors and matrices. For  $\mathbf{a}, \mathbf{b} \in \mathbb{I}^*\mathbb{R}$

$$\mathbf{dual}(\mathbf{dual}(\mathbf{a})) = \mathbf{a}, \quad \mathbf{dual}(\mathbf{a} \circ \mathbf{b}) = \mathbf{dual}(\mathbf{a}) \circ \mathbf{dual}(\mathbf{b}), \quad \circ \in \{+, -, \times, /\}.$$

The generalized interval arithmetic structure possesses group properties with respect to addition and multiplication operations. Thus,  $\mathbf{a} \ominus \mathbf{b} = \mathbf{a} - \mathbf{dual}(\mathbf{b})$ .

We will use the following sub-distributive property, proven in [1],

$$(5) \quad \mathbf{impr}(\mathbf{a})\mathbf{b} + \mathbf{ac} \subseteq \mathbf{a}(\mathbf{b} + \mathbf{c}),$$

wherein  $\mathbf{impr}(\mathbf{a}) := \{\mathbf{a} \text{ if } \underline{a} \geq \bar{a}, \mathbf{dual}(\mathbf{a}) \text{ if } \underline{a} \leq \bar{a}\}$  and  $\mathbf{a}, \mathbf{b}, \mathbf{c} \in \mathbb{I}^*\mathbb{R}$ .

Let  $\mathbb{F} \subset \mathbb{R}$  be the set of floating-point numbers and  $\mathbb{IF}, \mathbb{I}^*\mathbb{F}$  be the corresponding interval sets. Denote by  $\diamond$  the outward and by  $\circ$  the inward rounding,  $\diamond, \circ : \mathbb{I}^*\mathbb{R} \rightarrow \mathbb{I}^*\mathbb{F}$ . The following properties (cf. [1]) show that inner numerical approximations can be obtained at no additional cost only by outward directed rounding and the  $\mathbf{dual}$  operator in  $\mathbb{I}^*\mathbb{F}$ . For  $\mathbf{a}, \mathbf{b} \in \mathbb{I}^*\mathbb{F}$ ,  $\circ \in \{+, -, \times, /\}$ ,

$$(6) \quad \circ(\mathbf{a}) = \mathbf{dual}(\diamond(\mathbf{dual}\mathbf{a})); \quad \mathbf{a} \odot \mathbf{b} = \mathbf{dual}(\mathbf{dual}\mathbf{a} \diamond \mathbf{dual}\mathbf{b}).$$

**3. Inner estimation of linear  $\Sigma_{AE}^p$ .** We consider linear systems (1) and their parametric  $AE$ -solution sets which have linear shape. In this case the complete characterization of  $\Sigma_{AE}^p$  is given by a system of linear interval inclusions (4). We assume knowledge of an interior point  $\tilde{x}$  for  $\Sigma_{AE}^p$ . By Theorem 2.1, the relation

$$(7) \quad \sum_{k \in \mathcal{A}} \mathbf{p}_k(U_k \tilde{x} - v_k) \subseteq \text{int} \left( v_0 - U_0 \tilde{x} + \sum_{k \in \mathcal{E}} \mathbf{p}_k(v_k - U_k \tilde{x}) \right)$$

presents a sufficient condition for a point  $\tilde{x} \in \mathbb{R}^n$  to be in the interior of  $\Sigma_{AE}^p$ . It is an open problem how to find a suitable interior point  $\tilde{x}$ . A first natural trial point is the approximate midpoint solution. It can be found either by solving the quadratic system  $A(\check{\mathbf{p}})x = b(\check{\mathbf{p}})$  if  $m = n$ , or by solving the overdetermined system  $U(\check{\mathbf{p}})x = v(\check{\mathbf{p}})$  by least squares. It has to be verified by relation (7) that the chosen  $\tilde{x}$  is an interior point for  $\Sigma_{AE}^p$ .

For a given  $\tilde{x} \in \text{int } \Sigma_{AE}^p$ , we compute the maximal nonnegative number  $\eta$ , such that

$$(8) \quad \eta \left( U_0 \mathbf{e} + \sum_{k=1}^K \mathbf{p}_k(U_k \mathbf{e}) \right) \subseteq v_0 - U_0 \tilde{x} + \sum_{\mu \in \mathcal{E}} \mathbf{p}_\mu(v_\mu - U_\mu \tilde{x}) \ominus \sum_{\nu \in \mathcal{A}} \mathbf{p}_\nu(U_\nu \tilde{x} - v_\nu),$$

where  $\mathbf{e} = ([-1, 1], \dots, [-1, 1])^\top$ . Note that  $\mathbf{u} \ominus \mathbf{v} = \emptyset$  if  $\hat{\mathbf{v}} > \hat{\mathbf{u}}$ . However, for  $\tilde{x} \in \text{int } \Sigma_{AE}^p$ , (4), respectively (7), implies that the right-hand side of (8) will not be an empty set.

**Theorem 3.1.** *For  $\tilde{x} \in \text{int } \Sigma_{AE}^p$  and  $\eta \geq 0$ , such that (8) holds,*

$$\tilde{x} + \eta \mathbf{e} \subseteq \Sigma_{AE}^p.$$

**Proof.** For each  $x \in \tilde{x} + \eta \mathbf{e}$  we have

$$\begin{aligned} U_0 x + \sum_{\nu \in \mathcal{A}} \mathbf{p}_\nu(U_\nu x - v_\nu) &\subseteq U_0(\tilde{x} + \eta \mathbf{e}) + \sum_{\nu \in \mathcal{A}} \mathbf{p}_\nu(U_\nu(\tilde{x} + \eta \mathbf{e}) - v_\nu) \\ &\stackrel{\text{sub distr.}}{\subseteq} U_0 \tilde{x} + \eta U_0 \mathbf{e} + \sum_{\nu \in \mathcal{A}} \mathbf{p}_\nu(U_\nu \tilde{x} - v_\nu) + \eta \sum_{\nu \in \mathcal{A}} \mathbf{p}_\nu(U_\nu \mathbf{e}). \end{aligned}$$

By (8) the last expression to the right is contained in

$$\begin{aligned} U_0 \tilde{x} + \sum_{\nu \in \mathcal{A}} \mathbf{p}_\nu(U_\nu \tilde{x} - v_\nu) \ominus \eta \sum_{\mu \in \mathcal{E}} \mathbf{p}_\mu(U_\mu \mathbf{e}) + v_0 \\ - U_0 \tilde{x} + \sum_{\mu \in \mathcal{E}} \mathbf{p}_\mu(v_\mu - U_\mu \tilde{x}) \ominus \sum_{\nu \in \mathcal{A}} \mathbf{p}_\nu(U_\nu \tilde{x} - v_\nu), \end{aligned}$$

which due to  $\mathbf{a} \ominus \mathbf{a} = 0$  is equivalent to

$$\begin{aligned} v_0 + \sum_{\mu \in \mathcal{E}} \mathbf{p}_\mu(v_\mu - U_\mu \tilde{x}) \ominus \eta \sum_{\mu \in \mathcal{E}} \mathbf{p}_\mu(U_\mu \mathbf{e}) &\stackrel{(5)}{\subseteq} v_0 + \sum_{\mu \in \mathcal{E}} \mathbf{p}_\mu(v_\mu - U_\mu \tilde{x} - \eta \text{dual}(U_\mu \mathbf{e})) \\ &\subseteq v_0 + \sum_{\mu \in \mathcal{E}} \mathbf{p}_\mu(v_\mu - U_\mu(\tilde{x} + \eta \mathbf{e})) \\ &= v_0 + \sum_{\mu \in \mathcal{E}} \mathbf{p}_\mu(v_\mu - U_\mu x). \end{aligned}$$

The last inclusion is due to  $\text{impr}(\mathbf{a}) \subseteq \mathbf{a}$  and the equality is due to the initial assumption  $x \in \tilde{x} + \eta \mathbf{e}$ .

Thus  $U_0 x + \sum_{\nu \in \mathcal{A}} \mathbf{p}_\nu(U_\nu x - v_\nu) \subseteq v_0 + \sum_{\mu \in \mathcal{E}} \mathbf{p}_\mu(v_\mu - U_\mu x)$  for each  $x \in \tilde{x} + \eta \mathbf{e}$  implies  $\tilde{x} + \eta \mathbf{e} \subseteq \Sigma_{AE}^p$  by (4) of Theorem 2.1.  $\square$

In computational terms, from (8)

$$\eta = \min_{1 \leq i \leq q} \frac{\min\{|\underline{r}_i|, |\bar{r}_i|\}}{\bar{d}_i},$$

where  $\mathbf{r} := v_0 - U_0 \tilde{x} + \sum_{\mu \in \mathcal{E}} \mathbf{p}_\mu(v_\mu - U_\mu \tilde{x}) \ominus \sum_{\nu \in \mathcal{A}} \mathbf{p}_\nu(U_\nu \tilde{x} - v_\nu)$  and  $\mathbf{d} := U_0 \mathbf{e} + \sum_{k=1}^K \mathbf{p}_k(U_k \mathbf{e})$ .

The formula for inner interval subtraction implies

$$\underline{r} = v(\check{p}) - U(\check{p})\tilde{x} - t, \quad \bar{r} = v(\check{p}) - U(\check{p})\tilde{x} + t, \quad t = \sum_{\mu \in \mathcal{E}} \hat{p}_\mu |v_\mu - U_\mu \tilde{x}| - \sum_{\nu \in \mathcal{A}} \hat{p}_\nu |v_\nu - U_\nu \tilde{x}|.$$

With the notation  $U_k := (u_{k,ij}) \in \mathbb{R}^{q \times n}$  for  $k = 0, \dots, K$ , we have

$$\bar{d}_i = \sum_{j=1}^n |u_{0,ij}| + \sum_{k=1}^K \left( \sum_{j=1}^n |u_{k,ij}| \right) \max\{|\underline{p}_k|, |\bar{p}_k|\}, \quad \bar{d}_i \neq 0.$$

In the presence of round-off errors, the resulting inner estimation is valid if  $\mathbf{r}$  is rounded inward,  $\bar{d}_i$  are rounded upwards,  $\eta$  is rounded downwards, and  $\tilde{x} + \eta \mathbf{e}$  is rounded inward.

One can use a correct computer implementation of the arithmetic on proper and improper intervals, and the property (6) for a computer implementation of the parametric centred approach in order to provide safe bounds for the inner estimation. Based on (6), the following computational chain provides a guaranteed inner estimation  $\mathbf{z}$  of the parametric  $AE$ -solution set.

$$\begin{aligned} \mathbf{d} &= U_0 \mathbf{e} + \sum_{k=1}^K \mathbf{p}_k(U_k \mathbf{e}), \\ \mathbf{r} &= \text{dual} \left( v_0 - U_0 \tilde{x} + \sum_{\mu \in \mathcal{E}} \text{dual}(\mathbf{p}_\mu)(v_\mu - U_\mu \tilde{x}) - \sum_{\nu \in \mathcal{A}} \mathbf{p}_\nu(v_\nu - U_\nu \tilde{x}) \right), \\ \eta &= \min_{1 \leq i \leq q} \{ \min\{|\underline{r}_i|, |\bar{r}_i|\} / \bar{d}_i \}, \\ \mathbf{z} &= \text{dual}(\tilde{x} + \eta \text{dual}(\mathbf{e})), \end{aligned}$$

where all arithmetic operations (except the division on the third line, which is a floating-point division rounded to near) are outwardly rounded floating-point interval operations in  $\mathbb{I}^* \mathbb{F}$ . Similarly, verifying the condition (7) in floating-point

arithmetic we have to compute the left-hand side with outward rounding and the right-hand side with inward rounding.

The computational complexity of the centred approach is  $O(qn)$ . However, obtaining the explicit representation of the parametric  $AE$ -solution set will require some additional effort. The parametric centred method is implemented in the environment of *Mathematica*<sup>®</sup>. Another implementation environment suitable for large sparse data is the C-XSC class library, [3,11].

**Example 3.1.** We look for an inner estimation of the united solution set  $\Sigma_{uni}^g$  to the parametric system modelling a resistive electrical network<sup>1</sup>, which is considered for visualization in [10], Example 5.3. The explicit description of  $\Sigma_{uni}^g$  by Fourier–Motzkin-like elimination of the parameters leads to

$$U(g) = \begin{pmatrix} g_1 + g_6 & -g_6 & 0 & 0 & 0 \\ -g_6 & g_2 + g_6 + g_7 & -g_7 & 0 & 0 \\ 0 & -g_7 & g_3 + g_7 + g_8 & -g_8 & 0 \\ 0 & 0 & -g_8 & g_4 + g_8 + g_9 & -g_9 \\ 0 & 0 & 0 & -g_9 & g_5 + g_9 \\ -g_1 & -g_2 - g_7 & g_7 & 0 & 0 \\ g_6 & -g_2 - g_6 & -g_3 - g_8 & g_8 & 0 \\ g_1 & g_2 & g_3 + g_8 & -g_8 & 0 \\ 0 & g_7 & -g_3 - g_7 & -g_4 - g_9 & g_9 \\ -g_6 & g_2 + g_6 & g_3 & g_4 + g_9 & -g_9 \\ -g_1 & -g_2 & -g_3 - 2g_8 & g_4 + 2g_8 + g_9 & -g_9 \\ 0 & 0 & g_8 & -g_4 - g_8 & -g_5 \\ 0 & -g_7 & g_3 + g_7 & g_4 & g_5 \\ g_6 & -g_2 - g_6 & -g_3 & -g_4 - 2g_9 & g_5 + 2g_9 \\ g_1 & g_2 & g_3 + 2g_8 & -g_4 - 2g_8 - 2g_9 & g_5 + 2g_9 \end{pmatrix}, v(g) = \begin{pmatrix} 10 \\ 0 \\ 10 \\ 0 \\ 0 \\ -10 \\ -10 \\ 20 \\ -10 \\ 10 \\ -20 \\ 0 \\ 10 \\ -10 \\ 20 \end{pmatrix}.$$

Note that  $U_i(g) = A_i(g)$  and  $v_i(g) = b_i(g)$  for  $i = 1, \dots, 5$ . The parameters vary within  $g_i \in [0.99, 1.01]$ ,  $i = 1, \dots, 9$ , as in [10]. The approximate midpoint solution is an interior point for the parametric united solution set. Then, the parametric centred approach gives the following (rounded inward) guaranteed inner estimation

$$([7.084, 7.098], [4.175, 4.189], [5.448, 5.461], [2.175, 2.189], [1.084, 1.098])^\top.$$

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<sup>1</sup>the so-called Okumura’s problem [7].

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