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# BOUNDS FOR TERNARY EQUIDISTANT CONSTANT WEIGHT CODES

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In this paper we explore the problem of finding bounds for equidistant constant weight codes with  $2 \le w < n \le 10$ . Optimal ternary equidistant constant weight codes have been constructed by combinatorial and computer methods.

1. Introduction. Consider a finite set of q elements and containing a distinguished element "zero". The choice of a set does not matter in our context and we will use the set  $Z_q$  of integers modulo q. Let  $Z_q^n$  be the set of n-tuples (or vectors) over  $Z_q$  and  $Z_q^{n,w}$  be the set of n-tuples over  $Z_q$  of Hamming weight w.

A code is called *constant weight* if all the codewords have the same weight w. A code is called *equidistant* if all the distances between distinct codewords are d. Let  $B_q(n,d)$  denote the maximum number of codewords in an equidistant code over  $Z_q$  of length n and distance d (called an (n, M, d; q) equidistant code) and  $B_q(n, d, w)$  denote the maximum number of codewords in an equidistant constant weight code over  $Z_q$  of length n, distance d, and weight w (called an (n, M, d, w; q) equidistant constant weight code ECWC).

Equidistant codes have been investigated by a large number of authors, mainly as examples of designs and other combinatorial objects. Some works published on this topic are [5], [6], [8], [12]. Constant weight codes have been studied by many authors. For some references for the binary case, see Brouwer et al. [3], Agrell [1] and for the ternary case, see Bogdanova [2] and Svanström [11]. A few papers study codes which are both equidistant and of constant weight, for example [7], [10] and [4]. The same problem is considered in this paper.

2. Preliminaries. Some bounds for ECWC are given by the following theorems: Theorem 1 (Trivial values).

$$B_3(n,d,w) = 1 \text{ if } d > 2w,$$

$$B_3(n,d,n) = 2.$$

**Theorem 2** (the Johnson bounds for ECWC). The maximum number of codewords in a q-ary ECWC satisfy the inequalities:

$$B_q(n,d,w) \le \frac{n}{n-w} B_q(n-1,d,w),$$

$$B_q(n,d,w) \leq \frac{n(q-1)}{w} B_q(n-1,d,w-1).$$

The proof of the Theorem 2 is the same as the proof of Johnson bound for constantweight codes [11].

**Theorem 3** [4]. For k = 1, 2, ..., n, if  $P_k^2(w) > P_k(d) P_k(0)$ , then

$$B_q(n, d, w) \le \frac{P_k^2(0) - P_k(d) P_k(0)}{P_k^2(w) - P_k(d) P_k(0)}.$$

Here  $P_k(x)$  is the Krawtchouk polinomial defined by

$$P_{k}(x) = \sum_{i=0}^{k} {x \choose i} {n-x \choose k-i} (-1)^{i} (q-1)^{k-i}$$

and

$$P_k(0) = \binom{n}{k} (q-1)^k.$$

According to [7] there are no ECWC of order q, length q+1, distance q and weight q-1 which have more than  $\frac{(q^2+q)}{2}$  codewords, regardless whether q is even or odd.

#### 3. Some combinatorial bounds and constructions of ECWC.

**Proposition 4.** There exists a family of optimal ternary ECWC with parameters (n, 3, 3, 2; 3) for every integer  $n \ge 3$ .

**Proof.** Let u be a fixed codeword with length n and weight 2. Consider how many codewords are at distance exactly 3 from u we obtain that  $B_3(n, 3, 2) = 3$ .

**Proposition 5.** For w = 2, ..., n,  $P_n(w) \neq 0$ .

Proof.

$$P_n(w) = \sum_{i=0}^{n} {w \choose i} {n-w \choose n-i} (-1)^i (q-1)^{n-i}.$$

For the validity of binomial coefficients the following conditions must be satisfied:

$$\left|\begin{array}{c} i \le w \\ n - i \le n - w \end{array}\right.$$

Therefore i = w. Then  $P_n(w) = (-1)^w (q-1)^{n-w} \neq 0$ .

**Proposition 6.** For d = 3, w = 3, ..., n and k = n,

$$B_q(n,d,w) < (q-1)^3 + 1.$$

Proof.

$$B_q(n, d, w) \le \frac{P_k^2(0) - P_k(0) P_k(d)}{P_k^2(w) - P_k(0) P_k(d)}$$

For d=3 and k=n and using the same reasoning as in the Proposition 5 we obtain that

$$P_n(3) = \sum_{i=0}^{n} {3 \choose i} {n-3 \choose n-i} (-1)^i (q-1)^{n-i} = -(q-1)^{n-3}.$$

We have  $P_n(0) = \binom{n}{n} (q-1)^n = (q-1)^n$  and consequently

$$B_q(n,d,w) \le \frac{(q-1)^{2n} + (q-1)^{2n-3}}{P_k^2(w) + (q-1)^{2n-3}} < (q-1)^3 + 1.$$

**Corollary 7.** There exists a family of optimal ternary ECWC with parameters  $(4 + \lambda, 8, 3, 3 + t; 3)$  for every integer  $\lambda \geq 0$  and  $0 \leq t \leq n - 3$ .

**Proof.** From the Simplex code S, which has parameters (4,9,3;3) we construct ECWC  $C=S\setminus\{0\}$ . From the code C we construct a family of  $(4+\lambda,8,3,3+t;3)$  ECWC C' in the following way:

$$C' = \left\{ \left( \underbrace{00 \dots 0}_{\lambda - t} \underbrace{11 \dots 1}_{t}, c \right) | c \in C \right\},\,$$

where  $\lambda \geq 0$  and  $0 \leq t \leq n-3$ . Therefore  $B_3(n,3,w) \geq 8$ . For these parameters Proposition 6 gives  $B_3(n,3,w) < 9$ . So

$$B_3(n,3,w) = 8.$$

**4. Bounds for ECWC.** For codes of small size we apply combinatorial reasoning. For the rest of the values of M we use specificly developed, computer algorithms.

The best known upper and lower bounds (and exact values when these coincide) for ternary ECWC of length  $n \leq 10$  are displayed in Table 1. If in a certain position only one number occurs, then this number is the exact value of  $B_3(n,d,w)$  and the corresponding codes are optimal. If two numbers are given, then the right one is the best known upper bound for  $B_3(n,d,w)$ , received from Theorem 2 and Theorem 3. The left one is the best known lower bound for  $B_3(n,d,w)$ , received by our computer algorithm (exhaustive search), which is of exponential complexity.

All the numbers in column d=3 are obtained by Corollary 7.

Table 1. Bounds for optimal ternary ECWC

n	w	d=3	d = 4	d = 5	d = 6	d=7	d = 8	d = 9	d = 10
3	2	3							
4	2	3	2						
	3	8	2						
5	2	3	2						
	3	8	5	2					
	4	8	5	2					
6	2	3	3						
	3	8	5	4	2				
	4	8	6	4	3				
	5	8	6	3	2				
7	2	3	3						
	3	8	7	4	2				
	4	8	7	7	3	2			
	5	8	6	6	3	2			
	6	8	6	7	2	2			
8	2	3	4						
	3	8	7	4	2	_	_		
	4	8	7	7	5	2	2		
	5	8	8	7	8	3	2		
	6	8	6	7	8	2	2		
	7	8	8	8	4	2	2		
9	2	3	4		2				
	3	8	7	4	3		0		
	4	8	7	7	9	3	2	0	
	5	8	8	7	9	5	3	2	
	6 7	8	8	7	11 12	6	3	3	
	8	8	8	8	9	5 3	3 2	2 2	
10	2	3	5	0	9	3			
10	$\frac{2}{3}$	8	7	4	3				
	4	8	7	7	15	5	2		
	5	8	8	7	12 - 21	8 - 9	4	2	2
	6	8	8	7	14 - 20	8 - 15	5	3	2
	7	8	8	8	11 - 20	9 - 16	5	3	2
	8	8	8	8	12 - 21	10 - 11	5	2	2
	9	8	8	8	10	5	2	2	2
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## ГРАНИЦИ ЗА ТРОИЧНИ ЕКВИДИСТАНТНО КОНСТАНТНО ТЕГЛОВНИ КОДОВЕ

#### Галина Т. Богданова, Теодора А. Йоргова

Изследван е проблема за намиране на граници за троични еквидистантни константно тегловни кодове при  $2 \le w < n \le 10$ . Използвани са комбинаторни и компютърни методи за конструиране на оптимални троични еквидистантни константно тегловни кодове.